

TECHNICAL INFORMATION EXPLANATION

This technical information portion of your CP/M-85 documentation package is intended for users who wish to modify the CP/M Operating System or examine the internal components of the system.

Knowledge of this information is not necessary for most CP/M-85 users who only wish to use CP/M resident commands and to run utilities and application programs with the CP/M system.

This technical information consists of the following documentation items bound within three booklets:

- CP/M 2 System Interface: Chapter 5
- CP/M 2 Alteration: Chapter 6
- Appendix A: The MDS Basic I/O System (BIOS)
- Appendix B: A Skeletal CBIOS
- Appendix C: A Skeletal GETSYS/PUTSYS Program
- Appendix D: The MDS-800 Cold Start Loader for CP/M 2
- Appendix E: A Skeletal Cold Start Loader
- Appendix F: CP/M Disk Definition Library
- Appendix G: Blocking and Deblocking Algorithms.

This technical information was written by Digital Research Corporation, the original producers of CP/M Version 2. The source listings in the appendices are for illustration purposes only and do not correspond to the modules actually supplied with your system.

The CP/M-85 software products you have purchased are modifications of CP/M Version 2. In these modifications, the Basic Input/Output System (BIOS) has been customized for your hardware by Zenith Data Systems and Heath. Source listings for the actual modules used by your system are contained on the disks supplied with the system.

CP/M[®]
OPERATING SYSTEM
MANUAL

CP/M 2 System Interface
Chapter 5

 **DIGITAL RESEARCH[™]**
P.O. Box 579
Pacific Grove, California 93950

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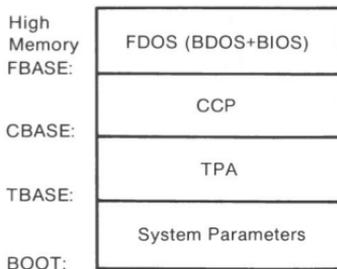
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CP/M 2 System Interface

5.1 Introduction

This chapter describes CP/M, release 2, system organization including the structure of memory and system entry points. The intention is to provide necessary information required to write programs that operate under CP/M and that use the peripheral and disk I/O facilities of the system.

CP/M is logically divided into four parts, called the Basic I/O System (BIOS), the Basic Disk Operating System (BDOS), the Console Command Processor (CCP), and the Transient Program Area (TPA). The BIOS is a hardware-dependent module that defines the exact low level interface with a particular computer system that is necessary for peripheral device I/O. Although a standard BIOS is supplied by Digital Research, explicit instructions are provided for field reconfiguration of the BIOS to match nearly any hardware environment (see Chapter 6). The BIOS and BDOS are logically combined into a single module with a common entry point and referred to as the FDOS. The CCP is a distinct program that uses the FDOS to provide a human-oriented interface with the information that is cataloged on the backup storage device. The TPA is an area of memory (i.e., the portion that is not used by the FDOS and CCP) where various nonresident operating system commands and user programs are executed. The lower portion of memory is reserved for system information and is detailed in later sections. Memory organization of the CP/M system is shown below.



The exact memory addresses corresponding to BOOT, TBASE, CBASE, and FBASE vary from version to version and are described fully in Chapter 6. All standard CP/M versions, however, assume BOOT = 0000H, which is the base of random access memory. The machine code found at location BOOT performs a system "warm start," which loads and initializes the programs and variables necessary to return control to the CCP. Thus, transient programs need only jump to location BOOT to return control to CP/M at the command level. Further, the standard versions assume TBASE = BOOT+0100H, which is normally location 0100H. The principal entry point to the FDOS is at location BOOT+0005H (normally 0005H) where a jump to FBASE is found. The address field at BOOT+0006H (normally 0006H) contains the value of FBASE and can be used to determine the size of available memory, assuming that the CCP is being overlaid by a transient program.

Transient programs are loaded into the TPA and executed as follows. The operator communicates with the CCP by typing command lines following each prompt. Each command line takes one of the forms:

```
command
command file1
command file1 file2
```

where "command" is either a built-in function such as DIR or TYPE or the name of a transient command or program. If the command is a built-in function of CP/M, it is executed immediately. Otherwise, the CCP searches the currently addressed disk for a file by the name

```
command.COM
```

If the file is found, it is assumed to be a memory image of a program that executes in the TPA and thus implicitly originates at TBASE in memory. The CCP loads the COM file from the disk into memory starting at TBASE and can extend up to CBASE.

If the command is followed by one or two file specifications, the CCP prepares one or two file control block (FCB) names in the system parameter area. These optional FCBs are in the form necessary to access files through the FDOS and are described in the next section.

The transient program receives control from the CCP and begins execution, using the I/O facilities of the FDOS. The transient program is "called" from the CCP. Thus, it can simply return to the CCP upon completion of its processing or can jump to BOOT to pass control back to CP/M. In the first case, the transient program must not use memory above CBASE, while in the latter case, memory up through FBASE-1 can be used.

The transient program can use the CP/M I/O facilities to communicate with the operator's console and peripheral devices, including the disk subsystem. The I/O system is accessed by passing a function number and an information address to CP/M through the FDOS entry point at BOOT+0005H. In the case of a disk read, for example, the transient program sends the number corresponding to a disk read, along with the address of an FCB to the CP/M FDOS. The FDOS, in turn, performs the operation and returns with either a disk read completion indication or an error number indicating that the disk read was unsuccessful.

5.2 Operating System Call Conventions

This section provides detailed information for performing direct operating system calls from user programs. Many of the functions listed below, however, are accessed more simply through the I/O macro library provided with the MAC macro assembler and listed in the Digital Research manual entitled, *MAC Macro Assembler: Language Manual and Applications Guide*.

CP/M facilities that are available for access by transient programs fall into two general categories: simple device I/O and disk file I/O. The simple device operations include:

- Read a Console Character
- Write a Console Character
- Read a Sequential Tape Character
- Write a Sequential Tape Character
- Write a List Device Character
- Get or Set I/O Status
- Print Console Buffer
- Read Console Buffer
- Interrogate Console Ready

The FDOS operations that perform disk I/O are

- Disk System Reset
- Drive Selection
- File Creation
- File Open
- File Close
- Directory Search
- File Delete
- File Rename
- Random or Sequential Read
- Random or Sequential Write
- Interrogate Available Disks
- Interrogate Selected Disk
- Set DMA Address
- Set/Reset File Indicators.

As mentioned above, access to the FDOS functions is accomplished by passing a function number and information address through the primary point at location BOOT+0005H. In general, the function number is passed in register C with the information address in the double byte pair DE. Single byte values are returned in register A, with double byte values returned in HL (a zero value is returned when the function number is out of range). For reasons of compatibility, register A = L and register B = H upon return in all cases. The user should note that the register passing conventions of CP/M agree with those of Intel's PL/M systems programming language. CP/M functions and their numbers are listed below.

0	System Reset	19	Delete File
1	Console Input	20	Read Sequential
2	Console Output	21	Write Sequential
3	Reader Input	22	Make File
4	Punch Output	23	Rename File
5	List Output	24	Return Login Vector
6	Direct Console I/O	25	Return Current Disk
7	Get I/O Byte	26	Set DMA Address
8	Set I/O Byte	27	Get Addr(Alloc)
9	Print String	28	Write Protect Disk
10	Read Console Buffer	29	Get R/O Vector
11	Get Console Status	30	Set File Attributes
12	Return Version Number	31	Get Addr(Disk Parms)
13	Reset Disk System	32	Set/Get User Code
14	Select Disk	33	Read Random
15	Open File	34	Write Random
16	Close File	35	Compute File Size
17	Search for First	36	Set Random Record
18	Search for Next	37	Reset Drive
		40	Write Random with Zero Fill

(Functions 28 and 32 should be avoided in application programs to maintain upward compatibility with CP/M.)

Upon entry to a transient program, the CCP leaves the stack pointer set to an eight-level stack area with the CCP return address pushed onto the stack, leaving seven levels before overflow occurs. Although this stack is usually not used by a transient program (i.e., most transients return to the CCP through a jump to location 0000H), it is sufficiently large to make CP/M system calls since the FDOS switches to a local stack at system entry. The assembly language program segment below, for example, reads characters continuously until an asterisk is encountered, at which time control returns to the CCP (assuming a standard CP/M system with BOOT = 0000H).

```

BDOS      EQU      0005H      ;STANDARD CP/M ENTRY
CONIN     EQU      1         ;CONSOLE INPUT FUNCTION
;
;
NEXTC:    ORG      0100H      ;BASE OF TPA
          MVI      C,CONIN    ;READ NEXT CHARACTER
          CALL     BDOS       ;RETURN CHARACTER IN <A>
          CPI      '*'        ;END OF PROCESSING?
          JNZ     NEXTC      ;LOOP IF NOT
          RET
          END                ;RETURN TO CCP

```

CP/M implements a named file structure on each disk, providing a logical organization that allows any particular file to contain any number of records from completely empty to the full capacity of the drive. Each drive is logically distinct with a disk directory and file data area. The disk file names are in three parts: the drive select code, the filename consisting of one to eight nonblank characters, and the filetype consisting of zero to three nonblank characters. The filetype names the generic category of a particular file, while the filename distinguishes individual files in each category. The filetypes listed below name a few generic categories that have been established, although they are somewhat arbitrary.

ASM	Assembler Source	PLI	PL/I Source File
PRN	Printer Listing	REL	Relocatable Module
HEX	Hex Machine Code	TEX	TEX Formatter Source
BAS	Basic Source File	BAK	ED Source Backup
INT	Intermediate Code	SYM	SID Symbol File
COM	Command File	\$\$\$	Temporary File

Source files are treated as a sequence of ASCII characters, where each "line" of the source file is followed by a carriage-return line-feed sequence (0DH followed by 0AH). Thus one 128-byte CP/M record could contain several lines of source text. The end of an ASCII file is denoted by a control-Z character (1AH) or a real end-of-file returned by the CP/M read operation. Control-Z characters embedded within machine code files (e.g., COM files) are ignored, however, and the end-of-file condition returned by CP/M is used to terminate read operations.

Files in CP/M can be thought of as a sequence of up to 65536 records of 128 bytes each, numbered from 0 through 65535, thus allowing a maximum of 8 megabytes per file. However, the user should note that although the records may be considered logically contiguous, they may not be physically contiguous in the disk data area. Internally, all files are divided into 16K byte segments called logical extents, so that counters are easily maintained as 8-bit values. The division into extents is discussed in the paragraphs that follow; however, they are not particularly significant for the programmer, since each extent is automatically accessed in both sequential and random access modes.

In the file operations starting with function number 15, DE usually addresses a file control block (FCB). Transient programs often use the default file control block area reserved by CP/M at location BOOT+005CH (normally 005CH) for simple file operations. The basic unit of file information is a 128-byte record used for all file operations; thus, a default location for disk I/O is provided by CP/M at location BOOT+0080H (normally 0080H), which is the initial default DMA address (see function 26). All directory operations take place in a reserved area that does not affect write buffers as was the case in release 1, with the exception of Search First and Search Next, where compatibility is required.

The FCB data area consists of a sequence of 33 bytes for sequential access and a series of 36 bytes in the case when the file is accessed randomly. The default FCB normally located at 005CH can be used for random access files, since the three bytes starting at BOOT+007DH are available for this purpose. The FCB format is shown with the following fields:

dr	f1	f2	/	/f8	t1	t2	t3	ex	s1	s2	rc	d0	/	/dn	cr	r0	r1	r2
00	01	02	...	08	09	10	11	12	13	14	15	16	...	31	32	33	34	35

where

dr	drive code (0-16) 0 => use default drive for file 1 => auto disk select drive A, 2 => auto disk select drive B, ... 16=> auto disk select drive P.
f1...f8	contain the file name in ASCII upper case, with high bit = 0
t1,t2,t3	contain the file type in ASCII upper case, with high bit = 0 t1', t2', and t3' denote the bit of these positions, t1' = 1 => Read/Only file, t2' = 1 => SYS file, no DIR list
ex	contains the current extent number, normally set to 00 by the user, but in range 0-31 during file I/O
s1	reserved for internal system use
s2	reserved for internal system use, set to zero on call to OPEN, MAKE, SEARCH
rc	record count for extent "ex," takes on values from 0-127
d0...dn	filled-in by CP/M, reserved for system use
cr	current record to read or write in a sequential file operation, normally set to zero by user
r0,r1,r2	optional random record number in the range 0-65535, with overflow to r2, r0, r1 constitute a 16-bit value with low byte r0, and high byte r1

Each file being accessed through CP/M must have a corresponding FCB, which provides the name and allocation information for all subsequent file operations. When accessing files, it is the programmer's responsibility to fill the lower 16 bytes of the FCB and initialize the cr field. Normally, bytes 1 through 11 are set to the ASCII character values for the file name and file type, while all other fields are zero.

FCBs are stored in a directory area of the disk, and are brought into central memory before the programmer proceeds with file operations (see the OPEN and MAKE functions). The memory copy of the FCB is updated as file operations take place and later recorded permanently on disk at the termination of the file operation (see the CLOSE command).

The CCP constructs the first 16 bytes of two optional FCBs for a transient by scanning the remainder of the line following the transient name, denoted by file1 and file2 in the prototype command line described above, with unspecified fields set to ASCII blanks. The first FCB is constructed at location BOOT+005CH and can be used as is for subsequent file operations. The second FCB occupies the d0 ... dn portion of the first FCB and must be moved to another area of memory before use. If, for example, the operator types

```
PROGNAME B:X.ZOT Y.ZAP
```

the file PROGNAME.COM is loaded into the TPA and the default FCB at BOOT+005CH is initialized to drive code 2, file name X, and file type ZOT. The second drive code takes the default value 0, which is placed at BOOT+006CH, with the file name Y placed into location BOOT+006DH and file type ZAP located 8 bytes later at BOOT+0075H. All remaining fields through cr are set to zero. The user should note again that it is the programmer's responsibility to move this second file name and type to another area, usually a separate file control block, before opening the file that begins at BOOT+005CH, because the open operation will overwrite the second name and type.

If no file names are specified in the original command, the fields beginning at BOOT+005DH and BOOT+006DH contain blanks. In all cases, the CCP translates lower case alphabetic to upper case to be consistent with the CP/M file naming conventions.

As an added convenience, the default buffer area at location BOOT+0080H is initialized to the command line tail typed by the operator following the program name. The first position contains the number of characters, with the characters themselves following the character count. Given the above command line, the area beginning at BOOT+0080H is initialized as follows:

BOOT+0080H:

```
+00 +01 +02 +03 +04 +05 +06 +07 +08 +09 +A +B +C +D +E  
E ' ' 'B' ' ' 'X' ' ' 'Z' 'O' 'T' ' ' 'Y' ' ' 'Z' 'A' 'P'
```

where the characters are translated to upper case ASCII with uninitialized memory following the last valid character. Again, it is the responsibility of the programmer to extract the information from this buffer before any file operations are performed, unless the default DMA address is explicitly changed.

Individual functions are described in detail in the pages that follow.

Function 0: System Reset

Entry Parameters:

Register C: 00H

The system reset function returns control to the CP/M operating system at the CCP level. The CCP reinitializes the disk subsystem by selecting and logging in disk drive A. This function has exactly the same effect as a jump to location BOOT.

Function 1: Console Input

Entry Parameters:

Register C: 01H

Returned Value:

Register A: ASCII Character

The console input function reads the next console character to register A. Graphic characters, along with carriage return, line feed, and back space (ctl-H) are echoed to the

console. Tab characters (ctl-I) move the cursor to the next tab stop. A check is made for start/stop scroll (ctl-S) and start/stop printer echo (ctl-P). The FDOS does not return to the calling program until a character has been typed, thus suspending execution if a character is not ready.

Function 2: Console Output

Entry Parameters:

Register C: 02H

Register E: ASCII Character

The ASCII character from register E is sent to the console device. As in function 1, tabs are expanded and checks are made for start/stop scroll and printer echo.

Function 3: Reader Input

Entry Parameters:

Register C: 03H

Returned Value:

Register A: ASCII Character

The Reader Input function reads the next character from the logical reader into register A (see the IOBYTE definition in Chapter 6). Control does not return until the character has been read.

Function 4: Punch Output

Entry Parameters:

Register C: 04H

Register E: ASCII Character

The Punch Output function sends the character from register E to the logical punch device.

Function 5: List Output

Entry Parameters:

Register C: 05H

Register E: ASCII Character

The List Output function sends the ASCII character in register E to the logical listing device.

Function 6: Direct Console I/O

Entry Parameters:

Register C: 06H
Register E: 0FFH (input) or
char (output)

Returned Value:

Register A: char or status

Direct console I/O is supported under CP/M for those specialized applications where basic console input and output are required. Use of this function should, in general, be avoided since it bypasses all of CP/M's normal control character functions (e.g., control-S and control-P). Programs that perform direct I/O through the BIOS under previous releases of CP/M, however, should be changed to use direct I/O under BDOS so that they can be fully supported under future releases of MP/M and CP/M.

Upon entry to function 6, register E either contains hexadecimal FF, denoting a console input request, or an ASCII character. If the input value is FF, function 6 returns A = 00 if no character is ready, otherwise A contains the next console input character.

If the input value in E is not FF, function 6 assumes that E contains a valid ASCII character that is sent to the console.

Function 6 must not be used in conjunction with other console I/O functions.

Function 7: Get I/O Byte

Entry Parameters:

Register C: 07H

Returned Value:

Register A: I/O Byte Value

The Get I/O Byte function returns the current value of IOBYTE in register A. See Chapter 6 for IOBYTE definition.

Function 8: Set I/O Byte

Entry Parameters:

Register C: 08H
Register E: I/O Byte Value

The Set I/O Byte function changes the IOBYTE value to that given in register E.

Function 9: Print String

Entry Parameters:

Register C: 09H
Registers DE: String Address

The Print String function sends the character string stored in memory at the location given by DE to the console device, until a \$ is encountered in the string. Tabs are expanded as in function 2, and checks are made for start/stop scroll and printer echo.

Function 10: Read Console Buffer

Entry Parameters:

Register C: 0AH
Registers DE: Buffer Address

Returned Value:

Console Characters in Buffer

The Read Buffer function reads a line of edited console input into a buffer addressed by registers DE. Console input is terminated when either input buffer overflows or a carriage return or line feed is typed. The Read Buffer takes the form:

```
DE: +0 +1 +2 +3 +4 +5 +6 +7 +8 ... +n
    |mx|nc |c1 |c2 |c3 |c4 |c5 |c6 |c7 |...|??|
```

where mx is the maximum number of characters that the buffer will hold (1 to 255) and nc is the number of characters read (set by FDOS upon return), followed by the characters read from the console. If nc < mx, then uninitialized positions follow the last character, denoted by ?? in the above figure. A number of control functions are recognized during line editing:

rub/del	removes and echoes the last character
ctl-C	reboots when at the beginning of line
ctl-E	causes physical end of line
ctl-H	backspaces one character position
ctl-J	(line feed) terminates input line
ctl-M	(return) terminates input line
ctl-R	retypes the current line after new line
ctl-U	removes current line
ctl-X	same as ctl-U.

The user should also note that certain functions that return the carriage to the leftmost position (e.g., ctl-X) do so only to the column position where the prompt ended (in earlier

releases, the carriage returned to the extreme left margin). This convention makes operator data input and line correction more legible.

Function 11: Get Console Status

Entry Parameters:

Register C: 0BH

Returned Value:

Register A: Console Status

The Console Status function checks to see if a character has been typed at the console. If a character is ready, the value 0FFH is returned in register A. Otherwise a 00H value is returned.

Function 12: Return Version Number

Entry Parameters:

Register C: 0CH

Returned Value:

Registers HL: Version Number

Function 12 provides information that allows version independent programming. A two-byte value is returned, with H = 00 designating the CP/M release (H = 01 for MP/M), and L = 00 for all releases previous to 2.0. CP/M 2.0 returns a hexadecimal 20 in register L, with subsequent version 2 releases in the hexadecimal range 21, 22, through 2F. Using function 12, for example, the user can write application programs that provide both sequential and random access functions.

Function 13: Reset Disk System

Entry Parameters:

Register C: 0DH

The Reset Disk Function is used to programmatically restore the file system to a reset state where all disks are set to read/write (see functions 28 and 29), only disk drive A is selected, and the default DMA address is reset to BOOT+0080H. This function can be used, for example, by an application program that requires a disk change without a system reboot.

Function 14: Select Disk

Entry Parameters:

Register C: 0EH

Register E: Selected Disk

The Select Disk function designates the disk drive named in register E as the default disk for subsequent file operations, with E = 0 for drive A, 1 for drive B, and so on through 15, corresponding to drive P in a full 16 drive system. The drive is placed in an on-line status, which activates its directory until the next cold start, warm start, or disk system reset operation. If the disk medium is changed while it is on-line, the drive automatically goes to a read/only status in a standard CP/M environment (see function 28). FCBs that specify drive code zero (dr = 00H) automatically reference the currently selected default drive. Drive code values between 1 and 16, however, ignore the selected default drive and directly reference drives A through P.

Function 15: Open File

Entry Parameters:

Register C: 0FH

Registers DE: FCB Address

Returned Value:

Register A: Directory Code

The Open File operation is used to activate a file that currently exists in the disk directory for the currently active user number. The FDOS scans the referenced disk directory for a match in positions 1 through 14 of the FCB referenced by DE (byte s1 is automatically zeroed), where an ASCII question mark (3FH) matches any directory character in any of these positions. Normally, no question marks are included, and bytes ex and s2 of the FCB are zero.

If a directory element is matched, the relevant directory information is copied into bytes d0 through dn of the FCB, thus allowing access to the files through subsequent read and write operations. The user should note that an existing file must not be accessed until a successful open operation is completed. Upon return, the open function returns a directory code with the value 0 through 3 if the open was successful or 0FFH (255 decimal) if the file cannot be found. If question marks occur in the FCB, the first matching FCB is activated. Note that the current record (cr) must be zeroed by the program if the file is to be accessed sequentially from the first record.

Function 16: Close File

Entry Parameters:

Register C: 10H
Registers DE: FCB Address

Returned Value:

Register A: Directory Code

The Close File function performs the inverse of the open file function. Given that the FCB addressed by DE has been previously activated through an open or make function (see functions 15 and 22), the close function permanently records the new FCB in the referenced disk directory. The FCB matching process for the close is identical to the open function. The directory code returned for a successful close operation is 0, 1, 2, or 3, while a 0FFH (255 decimal) is returned if the file name cannot be found in the directory. A file need not be closed if only read operations have taken place. If write operations have occurred, however, the close operation is necessary to record the new directory information permanently.

Function 17: Search for First

Entry Parameters

Register C: 11H
Registers DE: FCB Address

Returned Value:

Register A: Directory Code

Search First scans the directory for a match with the file given by the FCB addressed by DE. The value 255 (hexadecimal FF) is returned if the file is not found; otherwise, 0, 1, 2, or 3 is returned indicating the file is present. When the file is found, the current DMA address is filled with the record containing the directory entry, and the relative starting position is $A * 32$ (i.e., rotate the A register left 5 bits, or ADD A five times). Although not normally required for application programs, the directory information can be extracted from the buffer at this position.

An ASCII question mark (63 decimal, 3F hexadecimal) in any position from f1 through ex matches the corresponding field of any directory entry on the default or auto-selected disk drive. If the dr field contains an ASCII question mark, the auto disk select function is disabled and the default disk is searched, with the search function returning any matched entry, allocated or free, belonging to any user number. This latter function is not normally used by application programs, but it allows complete flexibility to scan all current directory values. If the dr field is not a question mark, the s2 byte is automatically zeroed.

Function 18: Search for Next

Entry Parameters:

Register C: 12H

Returned Value:

Register A: Directory Code

The Search Next function is similar to the Search First function, except that the directory scan continues from the last matched entry. Similar to function 17, function 18 returns the decimal value 255 in A when no more directory items match.

Function 19: Delete File

Entry Parameters:

Register C: 13H

Registers DE: FCB Address

Returned Value:

Register A: Directory Code

The Delete File function removes files that match the FCB addressed by DE. The filename and type may contain ambiguous references (i.e., question marks in various positions), but the drive select code cannot be ambiguous, as in the Search and Search Next functions.

Function 19 returns a decimal 255 if the referenced file or files cannot be found; otherwise, a value in the range 0 to 3 is returned.

Function 20: Read Sequential

Entry Parameters:

Register C: 14H

Registers DE: FCB Address

Returned Value:

Register A: Directory Code

Given that the FCB addressed by DE has been activated through an open or make function (numbers 15 and 22), the Read Sequential function reads the next 128-byte record from the file into memory at the current DMA address. The record is read from position cr of the extent, and the cr field is automatically incremented to the next record position. If the cr field overflows, the next logical extent is automatically opened and the cr field is reset to zero in preparation for the next read operation. The value 00H is returned in the A register if the read operation was successful, while a nonzero value is returned if no data exist at the next record position (e.g., end-of-file occurs).

Function 21: Write Sequential

Entry Parameters:

Register C: 15H
Registers DE: FCB Address

Returned Value:

Register A: Directory Code

Given that the FCB addressed by DE has been activated through an open or make function (numbers 15 and 22), the Write Sequential function writes the 128-byte data record at the current DMA address to the file named by the FCB. The record is placed at position cr of the file, and the cr field is automatically incremented to the next record position. If the cr field overflows, the next logical extent is automatically opened and the cr field is reset to zero in preparation for the next write operation. Write operations can take place into an existing file, in which case, newly written records overlay those that already exist in the file. Register A = 00H upon return from a successful write operation, while a nonzero value indicates an unsuccessful write caused by a full disk.

Function 22: Make File

Entry Parameters:

Register C: 16H
Registers DE: FCB Address

Returned Value:

Register A: Directory Code

The Make File operation is similar to the open file operation except that the FCB must name a file that does not exist in the currently referenced disk directory (i.e., the one named explicitly by a nonzero dr code or the default disk if dr is zero). The FDOS creates the file and initializes both the directory and main memory value to an empty file. The programmer must ensure that no duplicate file names occur, and a preceding delete operation is sufficient if there is any possibility of duplication. Upon return, register A = 0, 1, 2, or 3 if the operation was successful and 0FFH (255 decimal) if no more directory space is available. The make function has the side effect of activating the FCB and thus a subsequent open is not necessary.

Function 23: Rename File

Entry Parameters:

Register C: 17H
Registers DE: FCB Address

Returned Value:

Register A: Directory Code

The Rename function uses the FCB addressed by DE to change all occurrences of the file named in the first 16 bytes to the file named in the second 16 bytes. The drive code dr

at position 0 is used to select the drive, while the drive code for the new file name at position 16 of the FCB is assumed to be zero. Upon return, register A is set to a value between 0 and 3 if the rename was successful and 0FFH (255 decimal) if the first file name could not be found in the directory scan.

Function 24: Return Log-in Vector

Entry Parameters:

Register C: 18H

Returned Value:

Registers HL: Log-in Vector

The log-in vector value returned by CP/M is a 16-bit value in HL, where the least significant bit of L corresponds to the first drive A and the high order bit of H corresponds to the sixteenth drive, labeled P. A 0 bit indicates that the drive is not on-line, while a 1 bit marks a drive that is actively on-line as a result of an explicit disk drive selection or an implicit drive select caused by a file operation that specified a nonzero dr field. The user should note that compatibility is maintained with earlier releases, since registers A and L contain the same values upon return.

Function 25: Return Current Disk

Entry Parameters:

Register C: 19H

Returned Value:

Register A: Current Disk

Function 25 returns the currently selected default disk number in register A. The disk numbers range from 0 through 15 corresponding to drives A through P.

Function 26: Set DMA Address

Entry Parameters:

Register C: 1AH

Registers DE: DMA Address

DMA is an acronym for Direct Memory Address, which is often used in connection with disk controllers that directly access the memory of the mainframe computer to transfer data to and from the disk subsystem. Although many computer systems use non-DMA access (i.e., the data are transferred through programmed I/O operations), the DMA address has, in CP/M, come to mean the address at which the 128-byte data record resides before a disk write and after a disk read. Upon cold start, warm start, or disk

system reset, the DMA address is automatically set to BOOT+0080H. The Set DMA function, however, can be used to change this default value to address another area of memory where the data records reside. Thus, the DMA address becomes the value specified by DE until it is changed by a subsequent Set DMA function, cold start, warm start, or disk system reset.

Function 27: Get ADDR(Alloc)

Entry Parameters:

Register C: 1BH

Returned Value:

Registers HL: ALLOC Address

An allocation vector is maintained in main memory for each on-line disk drive. Various system programs use the information provided by the allocation vector to determine the amount of remaining storage (see the STAT program). Function 27 returns the base address of the allocation vector for the currently selected disk drive. However, the allocation information may be invalid if the selected disk has been marked read/only. Although this function is not normally used by application programs, additional details of the allocation vector are found in Chapter 6.

Function 28: Write Protect Disk

Entry Parameters:

Register C: 1CH

The disk write protect function provides temporary write protection for the currently selected disk. Any attempt to write to the disk before the next cold or warm start operation produces the message:

BDOS ERR on d: R/O

Function 29: Get Read/Only Vector

Entry Parameters:

Register C: 1DH

Returned Value:

Registers HL: R/O Vector Value

Function 29 returns a bit vector in register pair HL, which indicates drives that have the temporary read-only bit set. As in function 24, the least significant bit corresponds to drive A, while the most significant bit corresponds to drive P. The R/O bit is set either by an explicit call to function 28 or by the automatic software mechanisms within CP/M that detect changed disks.

Function 30: Set File Attributes

Entry Parameters:

Register C: 1EH
Registers DE: FCB Address

Returned Value:

Register A: Directory Code

The Set File Attributes function allows programmatic manipulation of permanent indicators attached to files. In particular, the R/O and System attributes (t1' and t2') can be set or reset. The DE pair addresses an unambiguous file name with the appropriate attributes set or reset. Function 30 searches for a match and changes the matched directory entry to contain the selected indicators. Indicators f1' through f4' are not currently used, but may be useful for applications programs, since they are not involved in the matching process during file open and close operations. Indicators f5' through f8' and t3' are reserved for future system expansion.

Function 31: Get ADDR(Disk Parms)

Entry Parameters:

Register C: 1FH

Returned Value:

Registers HL: DPB Address

The address of the BIOS resident disk parameter block is returned in HL as a result of this function call. This address can be used for either of two purposes. First, the disk parameter values can be extracted for display and space computation purposes, or transient programs can dynamically change the values of current disk parameters when the disk environment changes, if required. Normally, application programs will not require this facility.

Function 32: Set/Get User Code

Entry Parameters:

Register C: 20H
Register E: 0FFH (get) or
User Code (set)

Returned Value:

Register A: Current Code or
(no value)

An application program can change or interrogate the currently active user number by calling function 32. If register E = 0FFH, the value of the current user number is

returned in register A, where the value is in the range of 0 to 15. If register E is not 0FFH, the current user number is changed to the value of E (modulo 16).

Function 33: Read Random

Entry Parameters:

Register C: 21H
Registers DE: FCB Address

Returned Value:

Register A: Return Code

The Read Random function is similar to the sequential file read operation of previous releases, except that the read operation takes place at a particular record number, selected by the 24-bit value constructed from the 3-byte field following the FCB (byte positions r0 at 33, r1 at 34, and r2 at 35). The user should note that the sequence of 24 bits is stored with least significant byte first (r0), middle byte next (r1), and high byte last (r2). CP/M does not reference byte r2, except in computing the size of a file (function 35). Byte r2 must be zero, however, since a nonzero value indicates overflow past the end of file.

Thus, the r0, r1 byte pair is treated as a double-byte, or "word" value, which contains the record to read. This value ranges from 0 to 65535, providing access to any particular record of the 8-megabyte file. To process a file using random access, the base extent (extent 0) must first be opened. Although the base extent may or may not contain any allocated data, this ensures that the file is properly recorded in the directory and is visible in DIR requests. The selected record number is then stored in the random record field (r0, r1), and the BDOS is called to read the record. Upon return from the call, register A either contains an error code, as listed below, or the value 00, indicating the operation was successful. In the latter case, the current DMA address contains the randomly accessed record. The user should note that contrary to the sequential read operation, the record number is not advanced. Thus, subsequent random read operations continue to read the same record.

Upon each random read operation, the logical extent and current record values are automatically set. Thus, the file can be sequentially read or written, starting from the current randomly accessed position. However, the user should note that, in this case, the last randomly read record will be reread as one switches from random mode to sequential read and the last record will be rewritten as one switches to a sequential write operation. The user can, of course, simply advance the random record position following each random read or write to obtain the effect of a sequential I/O operation.

Error codes returned in register A following a random read are listed below.

01	reading unwritten data
02	(not returned in random mode)
03	cannot close current extent
04	seek to unwritten extent
05	(not returned in read mode)
06	seek past physical end of disk

Error codes 01 and 04 occur when a random read operation accesses a data block that has not been previously written or an extent that has not been created, which are equivalent conditions. Error code 03 does not normally occur under proper system

operation. If it does, it can be cleared by simply rereading or reopening extent zero as long as the disk is not physically write protected. Error code 06 occurs whenever byte r2 is nonzero under the current 2.0 release. Normally, nonzero return codes can be treated as missing data, with zero return codes indicating operation complete.

Function 34: Write Random

Entry Parameters:

Register C: 22H
Registers DE: FCB Address

Returned Value:

Register A: Return Code

The Write Random operation is initiated similarly to the Read Random call, except that data are written to the disk from the current DMA address. Further, if the disk extent or data block that is the target of the write has not yet been allocated, the allocation is performed before the write operation continues. As in the Read Random operation, the random record number is not changed as a result of the write. The logical extent number and current record positions of the file control block are set to correspond to the random record that is being written. Again, sequential read or write operations can begin following a random write, with the notation that the currently addressed record is either read or rewritten again as the sequential operation begins. The user can also simply advance the random record position following each write to get the effect of a sequential write operation. The user should note that, in particular, reading or writing the last record of an extent in random mode does not cause an automatic extent switch as it does in sequential mode.

The error codes returned by a random write are identical to the random read operation with the addition of error code 05, which indicates that a new extent cannot be created as a result of directory overflow.

Function 35: Compute File Size

Entry Parameters:

Register C: 23H
Registers DE: FCB Address

Returned Value:

Random Record Field Set

When computing the size of a file, the DE register pair addresses an FCB in random mode format (bytes r0, r1, and r2 are present). The FCB contains an unambiguous file name that is used in the directory scan. Upon return, the random record bytes contain the "virtual" file size, which is, in effect, the record address of the record following the end of the file. Following a call to function 35, if the high record byte r2 is 01, the file contains the maximum record count 65536. Otherwise, bytes r0 and r1 constitute a 16-bit value (r0 is the least significant byte, as before), which is the file size.

Data can be appended to the end of an existing file by simply calling function 35 to set the random record position to the end of file and then performing a sequence of random writes starting at the preset record address.

The virtual size of a file corresponds to the physical size when the file is written sequentially. If the file was created in random mode and "holes" exist in the allocation, the file may in fact contain fewer records than the size indicates. For example, if only the last record of an 8-megabyte file is written in random mode (i.e., record number 65535), the virtual size is 65536 records, although only one block of data is actually allocated.

Function 36: Set Random Record

Entry Parameters:

Register C: 24H

Registers DE: FCB Address

Returned Value:

Random Record Field Set

The Set Random Record function causes the BDOS automatically to produce the random record position from a file that has been read or written sequentially to a particular point. The function can be useful in two ways.

First, it is often necessary initially to read and scan a sequential file to extract the positions of various "key" fields. As each key is encountered, function 36 is called to compute the random record position for the data corresponding to this key. If the data unit size is 128 bytes, the resulting record position is placed into a table with the key for later retrieval. After scanning the entire file and tabulating the keys and their record numbers, the user can move instantly to a particular keyed record by performing a random read, using the corresponding random record number that was saved earlier. The scheme is easily generalized for variable record lengths, since the program need only store the buffer-relative byte position along with the key and record number to find the exact starting position of the keyed data at a later time.

A second use of function 36 occurs when switching from a sequential read or write over to random read or write. A file is sequentially accessed to a particular point in the file, function 36 is called, which sets the record number, and subsequent random read and write operations continue from the selected point in the file.

Function 37: Reset Drive

Entry Parameters:

Register C: 25H

Registers DE: Drive Vector

Returned Value:

Register A: 00H

The Reset Drive function allows resetting of specified drives. The passed parameter is a 16 bit vector of drives to be reset; the least significant bit is drive A:.

To maintain compatibility with MP/M, CP/M returns a zero value.

Function 40: Write Random With Zero Fill

Entry Parameters:

Register C: 28H
Registers DE: FCB Address

Returned Value:

Register A: Return Code

The Write Random With Zero Fill operation is similar to Function 34, with the exception that a previously unallocated block is filled with zeros before the data are written.

5.3 A Sample File-to-File Copy Program

The program shown below provides a relatively simple example of file operations. The program source file is created as COPY.ASM using the CP/M ED program and then assembled using ASM or MAC, resulting in a HEX file. The LOAD program is used to produce a COPY.COM file, which executes directly under the CCP. The program begins by setting the stack pointer to a local area and proceeds to move the second name from the default area at 006CH to a 33-byte file control block called DFCB. The DFCB is then prepared for file operations by clearing the current record field. At this point, the source and destination FCBs are ready for processing, since the SFCB at 005CH is properly set up by the CCP upon entry to the COPY program. That is, the first name is placed into the default FCB, with the proper fields zeroed, including the current record field at 007CH. The program continues by opening the source file, deleting any existing destination file, and creating the destination file. If all this is successful, the program loops at the label COPY until each record has been read from the source file and placed into the destination file. Upon completion of the data transfer, the destination file is closed and the program returns to the CCP command level by jumping to BOOT.

```

;
; sample file-to-file copy program
;
; at the ccp level, the command
;
; copy a:x.y b:u.v
;
; copies the file named x.y from drive
; a to a file named u.v. on drive b.
;
0000 = boot equ 0000h ; system reboot
0005 = bdos equ 0005h ; bdos entry point
005c = fcbl equ 005ch ; first file name
005c = sfcbl equ fcbl ; source fcb
006c = fcb2 equ 006ch ; second file name
0080 = dbuff equ 0080h ; default buffer
0100 = tpa equ 0100h ; beginning of tpa
;
0009 = printf equ 9 ; print buffer func#
000f = openf equ 15 ; open file func#
0010 = closef equ 16 ; close file func#
```



```

014e c33701                jmp copy    ; loop until eof
;
;
eofile:                    ; end of file, close destination
0151 11da01                lxi d,dfcb ; destination
0154 cd6e01                call close ; 255 if error
0157 21bb01                lxi h,wrprot ; ready message
015a 3c                    inr a      ; 255 becomes 00
015b cc6101                cz finis   ; shouldn't happen
;
;
; copy operation complete, end
015e 11cc01                lxi d,normal ; ready message
;
;
finis:                    ; write message given by de, reboot
0161 0e09                mvi c,printf
0163 cd0500                call bdos  ; write message
0166 c30000                jmp boot   ; reboot system
;
;
; system interface subroutines
; (all return directly from bdos)
;
;
0169 0e0f                open:     mvi c,openf
016b c30500                jmp bdos
;
;
016e 0e10                close:   mvi c,closef
0170 c30500                jmp bdos
;
;
0173 0e13                delete   mvi c,deletef
0175 c30500                jmp bdos
;
;
0178 0e14                read:    mvi c,readf
017a c30500                jmp bdos
;
;
017d 0e15                write:   mvi c,writef
017f c30500                jmp bdos
;
;
0182 0e16                make:    mvi c,makef
0184 c30500                jmp bdos
;
;
; console messages
0187 6e6f20f            nofile:  db 'no source file$'
0196 6e6f209            nodir:   db 'no directory spaces$'
01a9 6f7574f            space:   db 'out of data space$'
01bb 7772695            wrprot:  db 'write protected?$'
01cc 636f700            normal: db 'copy complete$'
;
;
; data areas
01da                    dfcb:    ds 33      ; destination fcb
01fa =                    dfcbcr   equ dfcb+32 ; current record
;
;
01fb                    ds 32      ; 16 level stack
;
stack:
021b                    end

```

The user should note that there are several simplifications in this particular program. First, there are no checks for invalid file names that could, for example, contain ambiguous references. This situation could be detected by scanning the 32-byte default area starting at location 005CH for ASCII question marks. A check should also be made to ensure that the file names have, in fact, been included (check locations 005DH and 006DH for nonblank ASCII characters). Finally, a check should be made to ensure that the source and destination file names are different. An improvement in speed could be obtained by buffering more data on each read operation. One could, for example, determine the size of memory by fetching FBASE from location 0006H and using the entire remaining portion of memory for a data buffer. In this case, the programmer simply resets the DMA address to the next successive 128-byte area before each read. Upon writing to the destination file, the DMA address is reset to the beginning of the buffer and incremented by 128 bytes to the end as each record is transferred to the destination file.

5.4 A Sample File Dump Utility

The file dump program shown below is slightly more complex than the simple copy program given in the previous section. The dump program reads an input file, specified in the CCP command line, and displays the content of each record in hexadecimal format at the console. Note that the dump program saves the CCP's stack upon entry, resets the stack to a local area, and restores the CCP's stack before returning directly to the CCP. Thus, the dump program does not perform a warm start at the end of processing.

```

; DUMP program reads input file and displays hex
; data
;
0100
0005 =      bdos      equ 0005h = ;bdos entry point
0001 =      cons      equ 1       ;read console
0002 =      typef     equ 2       ;type function
0009 =      printf    equ 9       ;buffer print entry
000b =      brkf      equ 11      ;break key function
;
000f =      openf     equ 15      ;file open
0014 =      readf     equ 20      ;read function
;
005c =      fcb       equ 5ch     ;file control block
;address
0080 =      buff      equ 80h     ;input disk buffer
;address
;
; non graphic characters
000d =      cr        equ 0dh     ;carriage return
000a =      lf        equ 0ah     ;line feed
;
; file control block definitions
005c =      fcbdn     equ fcb+0   ;disk name
005d =      fcbfn     equ fcb+1   ;file name
0065 =      fcbft     equ fcb+9   ;disk file type (3
;characters)
0068 =      fcbrl     equ fcb+12  ;file's current reel
;number
006b =      fcbrc     equ fcb+15  ;file's record count (0 to
;128)128)
007c =      fcbcr     equ fcb+32  ;current (next) record
;number (0

```

```

007d =          fcbln      equ fcb+33      ;fcb length
                ;
                ;
0100 210000     lxi      h,0
0103 39         dad      sp
                ;
0104 221502     ;          entry stack pointer in hl from the ccp
                ;          shld oldsp
                ;          set sp to local stack area (restored at
                ;          finis)
0107 315702     lxi      sp,stktp
                ;          read and print successive buffers
010a cdc101     call     setup      ;set up input file
010d feff       cpi      255      ;255 if file not present
010f c21b01     jnz      openok     ;skip if open is ok
                ;
                ;          file not there, give error message and
                ;          return
0112 11f301     lxi      d,opnmsg
0115 cd9c01     call     err
0118 c35101     jmp      finis      ;to return
                ;
openok:         ;open operation ok, set buffer index to
                ;end
011b 3e80       mvi      a,80h
011d 321302     sta      ibp      ;set buffer pointer to 80h
                ;          hl contains next address to print
0120 210000     lxi      h,0      ;start with 0000
                ;
gloop:         pushh     ;save line position
0123 e5         call     gnb
0124 cda201     pop      h      ;recall line position
0127 e1         jc      finis     ;carry set by gnb if end
0138 da5101     ;          ;file
012b 47         mov      b,a
                ;          print hex values
                ;          check for line fold
012c 7d         mov      a,l
012d e60f       ani      0fh      ;check low 4 bits
012f c24401     jnz      nonum
                ;          print line number
0132 cd7201     call     crlf
                ;
                ;          check for break key
0135 cd5901     call     break
                ;          accum lsb = 1 if character ready
0138 0f         rrc      ;into carry
0139 da5101     jc      finis     ;don't print any more
                ;
013c 7c         mov      a,h
013d cd8f01     call     phex
0140 7d         mov      a,l
0141 cd8f01     call     phex
                ;
nonum:         inx      h      ;to next line number
0144 23

```

```

0145 3e20          mvi a,' '
0147 cd6501       call pchar
014a 78           mov a,b
014b cd8f01       call phex
014e c32301       jmp gloop
;
;
finis:
;
;       end of dump, return to cco
;       (note that a jmp to 0000h reboots)
;
0151 cd7201       call crlf
0154 2a1502       lhid oldsp
0157 f9           sphl
;
;       stack pointer contains ccp's stack
;       location
0158 c9           ret          ;to the ccp
;
;
;
;       subroutines
;
break:
;check break key (actually any key will
;do)
0159 e5d5c5       push h! push d! push b; environment
; saved
015c 0e0b         mvi c,brkf
015e cd0500       call bdos
0161 c1d1e1       pop b! pop d! pop h; environment
restored
0164 c9           ret
;
pchar:
;print a character
0165 e5d5c5       push h! push d! push b; saved
0168 0e02         mvi c,typef
016a 5f           mov e,a
016b cd0500       call bdos
016e c1d1e1       pop b! pop d! pop h; restored
0171 c9           ret
;
;
crlf:
0172 3e0d         mvi a,cr
0174 cd6501       call pchar
0177 3e0a         mvi a,lf
0179 cd6501       call pchar
017c c9           ret
;
;
;
pnib:
;print nibble in reg a
017d e60f         ani 0fh          ;low 4 bits
017f fe0a         cpi 10
0181 d28901       jnc p10
;
;       less than or equal to 9
0184 c630         adi '0'
0186 c38b01       jmp prn
;
;
;
;       greater or equal to 10
0189 c637         p10: adi 'a' - 10

```

```

018b cd6501      prn:      call pchar
018e c9          ;          ret
;
;          phex:     ;print hex char in reg a
018f f5          ;          pushpsw
0190 0f          ;          rrc
0191 0f          ;          rrc
0192 0f          ;          rrc
0193 0f          ;          rrc
0194 cd7d01      call pnib      ;print nibble
0197 f1          pop psw
0198 cd7d01      call pnib
019b c9          ;          ret
;
;          err:     ;print error message
;          ;          d,e addresses message ending with "$"
019c 0e09        mvi c,printf  ;print buffer
;          ;          ;function
019e cd0500      call bdos
01a1 c9          ;          ret
;
;          ;
;          gnb:     ;get next byte
01a2 3a1302      lda ibp
01a5 fe80        cpi 80h
01a7 c2b301      jnz g0
;          ;          read another buffer
;          ;
;          ;
01aa cdce01      call disk
01ad b7          ora a          ;zero value if read ok
01ae cab301      jz g0         ;for another byte
;          ;          end of data, return with carry set for eof
01b1 37          stc
01b2 c9          ;          ret
;          ;
;          ;
;          g0:     ;read the byte at buff+reg a
01b3 5f          mov e,a       ;ls byte of buffer index
01b4 1600        mvi d,0      ;double precision
;          ;          ;index to de
01b6 3c          inr a        ;index=index+1
01b7 321302      sta ibp      ;back to memory
;          ;          pointer is incremented
;          ;          save the current file address
01ba 218000      lxi h,buff
01bd 19          dad d
;          ;          absolute character address is in hl
01be 7e          mov a,m
;          ;          byte is in the accumulator
01bf b7          ora a        ;reset carry bit
01c0 c9          ;          ret
;          ;
;          ;
;          setup:  ;set up file
;          ;          open the file for input
01c1 af          xra a        ;zero to accum

```


type data:.

The operator then responds by typing up to 127 characters, followed by a carriage return. RANDOM then writes the character string into the X.DAT file at record n. If the R command is issued, RANDOM reads record number n and displays the string value at the console. If the Q command is issued, the X.DAT file is closed, and the program returns to the CCP. In the interest of brevity, the only error message is

error, try again.

The program begins with an initialization section where the input file is opened or created, followed by a continuous loop at the label "ready" where the individual commands are interpreted. The default file control block at 005CH and the default buffer at 0080H are used in all disk operations. The utility subroutines then follow, which contain the principal input line processor, called "readc." This particular program shows the elements of random access processing, and can be used as the basis for further program development.

Sample Random Access Program for CP/M 2.0

```
0100                org      100h      ;base of tpa
;
0000 =             reboot    equ     0000h ;system reboot
0005 =             bdos     equ     0005h ;bdos entry point
;
0001 =             coninp   equ     1      ;console input function
0002 =             conout   equ     2      ;console output function
0009 =             pstring  equ     9      ;print string until '$'
000a =             rstring  equ     10     ;read console buffer
000c =             version  equ     12     ;return version number
000f =             openf    equ     15     ;file open function
0010 =             closef   equ     16     ;close function
0016 =             makef    equ     22     ;make file function
0021 =             readr    equ     33     ;read random
0022 =             writer   equ     34     ;write random
;
005c =             fcb      equ     005ch ;default file control
;block
007d =             ranrec   equ     fcb+33 ;random record position
007f =             ranovf   equ     fcb+35 ;high order (overflow)
;byte
0080 =             buff     equ     0080h ;buffer address
;
000d =             cr       equ     0dh    ;carriage return
000a =             lf       equ     0ah    ;line feed
;
```

Load SP, Set-Up File for Random Access

```
0100 31bc00                lxi      sp,stack
;
;
0103 0e0c                mvi      c,version
```

```

0105 cd0500      call    bdos
0108 fe20        cpi     20h      ;version 2.0 or better?
010a d21600      jnc     versok
;
;                bad version, message and go back
010d 111b00      lxi     d,badver
0110 cdda00      call    print
0113 c30000      jmp     reboot
;
;                versok:
;                correct version for random access
0116 0e0f        mvi     c,openf ;open default fcb
0118 115c00      lxi     d,fcbl
011b cd0500      call    bdos
011e 3c          inr     a        ;err 255 becomes zero
011f c23700      jnz     ready
;
;                cannot open file, so create it
0122 0e16        mvi     c,makef
0124 115c00      lxi     d,fcbl
0127 cd0500      call    bdos
012a 3c          inr     a        ;err 255 becomes zero
012b c23700      jnz     ready
;
;                cannot create file, directory full
012e 113a00      lxi     d,nospace
0131 cdda00      call    print
0134 c30000      jmp     reboot  ;back to ccp

```

Loop Back to Ready After Each Command

```

;
;                ready:
;                file is ready for processing
;
0137 cde500      call    readcom ;read next command
013a 227c00      shld   ranrec  ;store input record#
013d 217f00      lxi     h,ranovf
0140 3600        mvi     m,0    ;clear high byte if set
0142 fe51        cpi     'Q'    ;quit?
0144 c25600      jnz     notq
;
;                quit processing, close file
0147 0e10        mvi     c,closef
0149 115c00      lxi     d,fcbl
014c cd0500      call    bdos
014f 3c          inr     a        ;err 255 becomes 0
0150 cab900      jz     error   ;error message, retry
0153 c30000      jmp     reboot  ;back to ccp
;

```

End of Quit Command, Process Write

```
notq:
;       not the quit command, random write?
0156 fe57      cpi      'W'
0158 c28900    jnz      notw

;
;       this is a random write, fill buffer until cr
015b 114d00    lxi      d,datmsg
015e cdda00    call     print      ;data prompt
0161 0e7f      mvi      c,127      ;up to 127 characters
0163 218000    lxi      h,buff     ;destination
loop:      ;read next character to buff
0166 c5        push     b          ;save counter
0167 e5        push     h          ;next destination
0168 cdc200    call     getchr    ;character to a
016b e1        pop      h          ;restore counter
016c c1        pop      b          ;restore next to fill
016d fe0d      cpi      cr        ;end of line?
016f ca7800    jz       erloop

;       not end, store character
0172 77        mov      m,a
0173 23        inx     h          ;next to fill
0174 0d        dcr     c          ;counter goes down
0175 c26600    jnz     rloop     ;end of buffer?
erloop:
;       end of read loop, store 00
0178 3600     mvi     m,0

;
;       write the record to selected record number
017a 0e22     mvi     c,writer
017c 115c00   lxi     d,fcbl
017f cd0500   call    bdos
0182 b7       ora     a          ;error code zero?
0183 c2b900   jnz    error      ;message if not
0186 c33700   jmp    ready      ;for another record
;
```

End of Write Command, Process Read

```
notw:
;       not a write command, read record?
0189 fe52      cpi      'R'
018b c2b900    jnz     error     ;skip if not

;
;       read random record
018e 0e21     mvi     c,readr
0190 115c00   lxi     d,fcbl
0193 cd0500   call    bdos
0196 b7       ora     a          ;return code 00?
0197 c2b900   jnz    error

;
;       read was successful, write to console
```

019a cdcf00	call	crlf	;new line
019d 0e80	mvi	c,128	;max 128 characters
019f 218000	lxi	h,buf	;next to get
wloop:			
01a2 7e	mov	a,m	;next character
01a3 23	inx	h	;next to get
01a4 e67f	ani	7fh	;mask parity
01a6 ca3700	jz	ready	;for another command
			;if 00
01a9 c5	push	b	;save counter
01aa e5	push	h	;save next to get
01ab fe20	cpi	' '	;graphic?
01ad d4c800	cnc	putchr	;skip output if not
01b0 e1	pop	h	
01b1 c1	pop	b	
01b2 0d	dcr	c	;count=count-1
01b3 c2a200	jnz	wloop	
01b6 c33700	jmp	ready	

End of Read Command, All Errors End Up Here

```

;
error:
01b9 115900      lxi    d,errmsg
01bc cdda00      call   print
01bf c33700      jmp    ready
;

```

Utility Subroutines for Console I/O

```

getchr:
;read next console character to a
mvi    c,coninp
call   bdos
ret

;
putchr:
;write character from a to console
mvi    c,conout
mov    e,a      ;character to send
call   bdos    ;send character
ret

;
crlf:
;send carriage return line feed
mvi    a,cr     ;carriage return
call   putchr
mvi    a,lf     ;line feed
call   putchr
ret

```

```

;
print:
;print the buffer addressed by de until $
push    d
01da d5      call    crlf
01db cdcf00  pop     d           ;new line
01de d1      mvi    c,pstring
01df 0e09   call   bdos        ;print the string
01e1 cd0500  ret
;
readcom:
;read the next command line to the conbuf
01e5 116b00  lxi    d,prompt
01e8 cdda00  call   print       ;command?
01eb 0e0a   mvi    c,rstring
01ed 117a00  lxi    d,conbuf
01f0 cd0500  call   bdos        ;read command line
; command line is present, scan it
01f3 210000  lxi    h,0         ;start with 0000
01f6 117c00  lxi    d,conlin   ;command line
01f9 1a      readc: ldax   d           ;next command
;character
01fa 13      inx    d           ;to next command
;position
01fb b7      ora    a           ;cannot be end of
;command
01fc c8      rz
; not zero, numeric?
01fd d630   sui    '0'
01ff fe0a   cpi    10         ;carry if numeric
0201 d21300  jnc    endr
; add-in next digit
0204 29     dad    h           ;*2
0205 4d     mov    c,l
0206 44     mov    b,h       ;bc = value * 2
0207 29     dad    h           ;*4
0208 29     dad    h           ;*8
0209 09     dad    b         ;*2 + *8 = *10
020a 85     add    l         ;+digit
020b 6f     mov    l,a
020c d2f900  jnc    readc     ;for another char
020f 24     inr    h         ;overflow
0210 c3f900  jmp    readc     ;for another char
endr:
; end of read, restore value in a
0213 c630   adi    '0'       ;command
0215 fe61   cpi    'a'       ;translate case?
0217 d8     rc
; lower case, mask lower case bits
0218 e65f   ani    101$1111b
021a c9     ret
;

```

String Data Area for Console Messages

```

badver:      db      'sorry, you need cp/m version 2$'
nospace:    db      'no directory space$'
datmsg:     db      'type data: $'
errmsg:     db      'error, try again.$'
prompt:     db      'next command? $'
;
    
```

Fixed and Variable Data Area

```

027a 21      conbuf: db      conlen ;length of console buffer
027b         consiz: ds      1      ;resulting size after read
027c         conlin: ds     32     ;length 32 buffer
0021 =       conlen equ     $-consiz
;
029c         ds      32      ;16 level stack
02bc        stack: end
    
```

Again, major improvements could be made to this particular program to enhance its operation. In fact, with some work, this program could evolve into a simple data base management system. One could, for example, assume a standard record size of 128 bytes, consisting of arbitrary fields within the record. A program, called GETKEY, could be developed that first reads a sequential file and extracts a specific field defined by the operator. For example, the command

```
GETKEY NAMES.DAT LASTNAME 10 20
```

would cause GETKEY to read the data base file NAMES.DAT and extract the "LASTNAME" field from each record, starting in position 10 and ending at character 20. GETKEY builds a table in memory consisting of each particular LASTNAME field, along with its 16-bit record number location within the file. The GETKEY program then sorts this list and writes a new file, called LASTNAME.KEY, which is an alphabetical list of LASTNAME fields with their corresponding record numbers. (This list is called an *inverted index* in information retrieval parlance.)

If the programmer were to rename the program shown above as QUERY and message it so that it reads a sorted key file into memory, the command line might appear as

```
QUERY NAMES.DAT LASTNAME.KEY.
```

Instead of reading a number, the QUERY program reads an alphanumeric string that is a particular key to find in the NAMES.DAT data base. Since the LASTNAME.KEY list is sorted, one can find a particular entry rapidly by performing a "binary search," similar to looking up a name in the telephone book. That is, starting at both ends of the list, one examines the entry halfway in between and, if not matched, splits either the upper half or

the lower half for the next search. The user will quickly reach the item he or she is looking for and find the corresponding record number. The user should fetch and display this record at the console, just as was done in the program shown above.

With some more work, the user can allow a fixed grouping size that differs from the 128-byte record shown above. This is accomplished by keeping track of the record number as well as the byte offset within the record. Knowing the group size, one randomly accesses the record containing the proper group, offset to the beginning of the group within the record read sequentially until the group size has been exhausted.

Finally, one can improve QUERY considerably by allowing boolean expressions, which compute the set of records that satisfy several relationships, such as a LASTNAME between HARDY and LAUREL and an AGE lower than 45. Display all the records that fit this description. Finally, if the user's lists are getting too big to fit into memory, he or she should randomly access key files from the disk as well.

5.6 System Function Summary

FUNCTION NUMBER	FUNCTION NAME	INPUT	OUTPUT
Decimal	Hex		
0	0 System Reset	C = 00H	none
1	1 Console Input	C = 01H	A = ASCII char
2	2 Console Output	E = char	none
3	3 Reader Input		A = ASCII char
4	4 Punch Output	E = char	none
5	5 List Output	E = char	none
6	6 Direct Console I/O	C = 06H E = OFFH (input) or OFEH (status) or char (output)	A = char or status (no value)
7	7 Get I/O Byte	none	A = I/O Byte Value
8	8 Set I/O Byte	E = I/O Byte	none
9	9 Print String	DE = Buffer Address	none
10	A Read Console Buffer	DE = Buffer	Console Characters in Buffer
11	B Get Console Status	none	A = 00/non zero
12	C Return Version Number	none	HL: Version Number
13	D Reset Disk System	none	none
14	E Select Disk	E = Disk Number	none
15	F Open File	DE = FCB Address	FF if not found
16	10 Close File	DE = FCB Address	FF if not found

CP/M[®]
OPERATING SYSTEM
MANUAL

CP/M 2 Alteration
Chapter 6

 **DIGITAL RESEARCH[™]**
P.O. Box 579
Pacific Grove, California 93950

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CP/M 2 Alteration

6.1 Introduction

The standard CP/M system assumes operation on an Intel MDS-800 microcomputer development system, but is designed so the user can alter a specific set of subroutines that define the hardware operating environment.

Although standard CP/M 2 is configured for single density floppy disks, field-alteration features allow adaptation to a wide variety of disk subsystems from single drive minidisks through high-capacity, "hard disk" systems. To simplify the following adaptation process, it is assumed that CP/M 2 will first be configured for single density floppy disks where minimal editing and debugging tools are available. If an earlier version of CP/M is available, the customizing process is eased considerably. In this latter case, the user may wish to review the system generation process and skip to later sections that discuss system alteration for nonstandard disk systems.

To achieve device independence, CP/M is separated into three distinct modules:

BIOS	basic I/O system, which is environment dependent
BDOS	basic disk operating system, which is not dependent upon the hardware configuration
CCP	the console command processor, which uses the BDOS

Of these modules, only the BIOS is dependent upon the particular hardware. That is, the user can "patch" the distribution version of CP/M to provide a new BIOS that provides a customized interface between the remaining CP/M modules and the user's own hardware system. This document provides a step-by-step procedure for patching a new BIOS into CP/M.

All disk-dependent portions of CP/M 2 are placed into a BIOS, a resident "disk parameter block," which is either hand coded or produced automatically using the disk definition macro library provided with CP/M 2. The end user need only specify the maximum number of active disks, the starting and ending sector numbers, the data allocation size, the maximum extent of the logical disk, directory size information, and reserved track values. The macros use this information to generate the appropriate tables and table references for use during CP/M 2 operation. Deblocking information is provided, which

aids in assembly or disassembly of sector sizes that are multiples of the fundamental 128 byte data unit, and the system alteration manual includes general purpose subroutines that use the deblocking information to take advantage of larger sector sizes. Use of these subroutines, together with the table-drive data access algorithms, makes CP/M 2 a universal data management system.

File expansion is achieved by providing up to 512 logical file extents, where each logical extent contains 16K bytes of data. CP/M 2 is structured, however, so that as much as 128K bytes of data are addressed by a single physical extent (corresponding to a single directory entry) maintaining compatibility with previous versions while taking advantage of directory space.

If CP/M is being tailored to a computer system for the first time, the new BIOS requires some simple software development and testing. The standard BIOS is listed in Appendix A and can be used as a model for the customized package. A skeletal version of the BIOS given in Appendix B can serve as the basis for a modified BIOS. In addition to the BIOS, the user must write a simple memory loader, called GETSYS, that brings the operating system into memory. To patch the new BIOS into CP/M, the user must write the reverse of GETSYS, called PUTSYS, which places an altered version of CP/M back onto the diskette. PUTSYS can be derived from GETSYS by changing the disk read commands into disk write commands. Sample skeletal GETSYS and PUTSYS programs are described in Section 6.4 and listed in Appendix C. To make the CP/M system load automatically, the user must also supply a cold start loader, similar to the one provided with CP/M (listed in Appendices A and D). A skeletal form of a cold start loader is given in Appendix E, which serves as a model for the loader.

6.2 First Level System Regeneration

The procedure to patch the CP/M system is given below. Address references in each step are shown with "H" denoting the hexadecimal radix, and are given for a 20K CP/M system. For larger CP/M systems, a "bias" is added to each address that is shown with a "+b" following it, where b is equal to the memory size—20K. Values for b in various standard memory sizes are

24K:	$b = 24K - 20K = 4K = 1000H$
32K:	$b = 32K - 20K = 12K = 3000H$
40K:	$b = 40K - 20K = 20K = 5000H$
48K:	$b = 48K - 20K = 28K = 7000H$
56K:	$b = 56K - 20K = 36K = 9000H$
62K:	$b = 62K - 20K = 42K = A800H$
64K:	$b = 64K - 20K = 44K = B000H$

It should be noted that the standard distribution version of CP/M is set for operation within a 20K memory system. Therefore, the user must first bring up the 20K CP/M system, then configure it for actual memory size (the user should see Section 6.3).

The user should:

1. Read Section 6.4 and write a GETSYS program that reads the first two tracks of a diskette into memory. The program from the diskette must be loaded starting at location 3380H. GETSYS is coded to start at location 100H (base of the TPA), as shown in Appendix C.

2. Test the GETSYS program by reading a blank diskette into memory and check to see that the data have been read properly and that the diskette has not been altered in any way by the GETSYS program.

3. Run the GETSYS program using an initialized CP/M diskette to see if GETSYS loads CP/M starting at 3380H (the operating system actually starts 128 bytes later at 3400H).

4. Read Section 6.4 and write the PUTSYS program. This writes memory starting at 3380H back onto the first two tracks of the diskette. The PUTSYS program should be located at 200H, as shown in Appendix C.

5. Test the PUTSYS program using a blank, uninitialized diskette by writing a portion of memory to the first two tracks; clear memory and read it back using GETSYS. Test PUTSYS completely, since this program will be used to alter CP/M on disk.

6. Study Sections 6.5, 6.6, and 6.7 along with the distribution version of the BIOS given in Appendix A and write a simple version that performs a similar function for the customized environment. Use the program given in Appendix B as a model. Call this new BIOS by the name CBIOS (customized BIOS). Implement only the primitive disk operations on a single drive and simple console input/output functions in this phase.

7. Test CBIOS completely to ensure that it properly performs console character I/O and disk reads and writes. Be careful to ensure that no disk write operations occur during read operations and check that the proper track and sectors are addressed on all reads and writes. Failure to make these checks may cause destruction of the initialized CP/M system after it is patched.

8. Referring to the table in Section 6.5, note that the BIOS is placed between locations 4A00H and 4FFFH. Read the CP/M system using GETSYS and replace the BIOS segment by the CBIOS developed in step 6 and tested in step 7. This replacement is done in memory.

9. Use PUTSYS to place the patched memory image of CP/M onto the first two tracks of a blank diskette for testing.

10. Use GETSYS to bring the copied memory image from the test diskette back into memory at 3380H and check to ensure that it has loaded back properly (clear memory, if possible, before the load). Upon successful load, branch to the cold start code at location 4A00H. The cold start routine will initialize page zero, then jump to the CCP at location 3400H, which will call the BDOS, which will call the CBIOS. The CBIOS will be asked by the CCP to read sixteen sectors on track 2, and CP/M will type "A>", the system prompt.

If difficulties are encountered, use whatever debug facilities are available to trace and breakpoint the CBIOS.

11. Upon completion of step 10, CP/M has prompted the console for a command input. Test the disk write operation by typing

SAVE 1 X.COM

(All commands must be followed by a carriage return.) CP/M responds with another prompt (after several disk accesses)

A>

If it does not, debug the disk write functions and retry.

12. Test the directory command by typing

DIR

CP/M responds with

A: X COM

13. Test the erase command by typing

ERA X.COM

CP/M responds with the A prompt. This is now an operational system that only requires a bootstrap loader to function completely.

14. Write a bootstrap loader that is similar to GETSYS and place it on track 0, sector 1 using PUTSYS (again using the test diskette, not the distribution diskette). See Sections 6.5 and 6.8 for more information on the bootstrap operation.

15. Retest the new test diskette with the bootstrap loader installed by executing steps 11, 12, and 13. Upon completion of these tests, type a control-C (control and C keys simultaneously). The system executes a "warm start" that reboots the system, and types the A prompt.

16. At this point, there is probably a good version of the customized CP/M system on the test diskette. Use GETSYS to load CP/M from the test diskette. Remove the test diskette, place the distribution diskette (or a legal copy) into the drive, and use PUTSYS to replace the distribution version with the customized version. The user should not make this replacement if unsure of the patch because this step destroys the system that was obtained from Digital Research.

17. Load the modified CP/M system and test it by typing

DIR

CP/M responds with a list of files that are provided on the initialized diskette. One file is the memory image for the debugger

DDT.COM

Note that from now on, it is important always to reboot the CP/M system (ctl-C is sufficient) when the diskette is removed and replaced by another diskette, unless the new diskette is to be read only.

18. Load and test the debugger by typing

DDT

(See Chapter 4 for operating procedures.)

19. Before making further CBIOS modifications, practice using the editor (see Chapter 2), and assembler (see Chapter 3). Recode and test the GETSYS, PUTSYS, and CBIOS programs using ED, ASM, and DDT. Code and test a COPY program that does a sector-to-sector copy from one diskette to another to obtain back-up copies of the original diskette. (Read the CP/M Licensing Agreement specifying legal responsibilities when copying the CP/M system.) Place the copyright notice

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on each copy that is made with the COPY program.

20. Modify the CBIOS to include the extra functions for punches, readers, and sign-on messages, and add the facilities for additional disk drives, if desired. These changes can be made with the GETSYS and PUTSYS programs or by referring to the regeneration process in Section 6.3.

The user should now have a good copy of the customized CP/M system. Although the CBIOS portion of CP/M belongs to the user, the modified version cannot be legally copied for anyone else's use.

It should be noted that the system remains file-compatible with all other CP/M systems (assuming media compatibility), which allows transfer of nonproprietary software between CP/M users.

6.3 Second Level System Generation

Once the system is running, the user will want to configure CP/M for the desired memory size. Usually a memory image is first produced with the "MOVCPM" program (system relocater) and then placed into a named disk file. The disk file can then be loaded, examined, patched, and replaced using the debugger and the system generation program. (The user should refer to Chapter 1.)

The CBIOS and BOOT are modified using ED and assembled using ASM, producing files called CBIOS.HEX and BOOT.HEX, which contain the code for CBIOS and BOOT in Intel hex format.

To get the memory image of CP/M into the TPA configured for the desired memory size, the user should type the command

```
MOVCPM xx *
```

where xx is the memory size in decimal K bytes (e.g., 32 for 32K). The response will be

```
CONSTRUCTING xxK CP/M VERS 2.0  
READY FOR "SYSGEN" OR  
"SAVE 34 CPMxx.COM"
```

An image of CP/M in the TPA is configured for the requested memory size. The memory image is at location 0900H through 227FH (i.e., the BOOT is at 0900H, the CCP is at 980H, the BDOS starts at 1180H, and the BIOS is at 1F80H). The user should note that the memory image has the standard MDS-800 BIOS and BOOT on it. It is now necessary to save the memory image in a file so that the user can patch the CBIOS and CBOOT into it:

```
SAVE 34 CPMxx.COM
```

Upon completion of the read, the user should reexamine the area where the CBIOS has been loaded (use an "L1F80" command) to ensure that it was loaded properly. When satisfied that the change has been made, the user should return from DDT using a control-C or, "GO" command.

SYSGEN is used to replace the patched memory image back onto a diskette (the user should utilize a test diskette until sure of the patch), as shown in the following interaction:

SYSGEN	Start the SYSGEN program
SYSGEN VERSION 2.0	Sign-on message from SYSGEN
SOURCE DRIVE NAME (OR RETURN TO SKIP)	Respond with a carriage return to skip the CP/M read operation since the system is already in memory
DESTINATION DRIVE NAME (OR RETURN TO REBOOT)	Respond with "B" to write the new system to the diskette in drive B
DESTINATION ON B, THEN TYPE RETURN	Place a scratch diskette in drive B, then type return.
FUNCTION COMPLETE DESTINATION DRIVE NAME (OR RETURN TO REBOOT)	

The user should place the scratch diskette in drive A and then perform a cold start to bring up the newly configured CP/M system.

The new CP/M system is then tested and the Digital Research copyright notice is placed on the diskette, as specified in the Licensing Agreement:

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6.4 Sample GETSYS and PUTSYS Programs

The following program provides a framework for the GETSYS and PUTSYS programs referenced in Sections 6.1 and 6.2. The READSEC and WRITESEC subroutines must be inserted by the user to read and write the specific sectors.

```

; GETSYS PROGRAM — READ TRACKS 0 AND 1 TO MEMORY AT 3380H
; REGISTER USE
;
; A (SCRATCH REGISTER)
; B TRACK COUNT (0, 1)
; C SECTOR COUNT (1,2, . . .,26)
; DE (SCRATCH REGISTER PAIR)
; HL LOAD ADDRESS
; SP SET TO STACK ADDRESS
;
START: LXI SP,3380H ;SET STACK POINTER TO SCRATCH
;AREA
LXI H, 3380H ;SET BASE LOAD ADDRESS
MVI B, 0 ;START WITH TRACK 0

```

```

RDTRK:                                ;READ NEXT TRACK (INITIALLY 0)
MVI C,1                                ;READ STARTING WITH SECTOR 1

RDSEC:                                ;READ NEXT SECTOR
CALL READSEC                            ;USER-SUPPLIED SUBROUTINE
LXI D,128                               ;MOVE LOAD ADDRESS TO NEXT 1/2
                                           ;PAGE
DAD D                                  ;HL = HL + 128
INR C                                  ;SECTOR = SECTOR + 1
MOV A,C                                ;CHECK FOR END OF TRACK
CPI 27
JC RDSEC                                ;CARRY GENERATED IF SECTOR < 27

;
; ARRIVE HERE AT END OF TRACK, MOVE TO NEXT TRACK
INR B
MOV A,B                                ;TEST FOR LAST TRACK
CPI 2
JC RDTRK                                ;CARRY GENERATED IF TRACK < 2

;
; ARRIVE HERE AT END OF LOAD, HALT FOR NOW
HLT

;
; USER-SUPPLIED SUBROUTINE TO READ THE DISK
READSEC:
; ENTER WITH TRACK NUMBER IN REGISTER B,
; SECTOR NUMBER IN REGISTER C, AND
; ADDRESS TO FILL IN HL
;
;
; PUSH B                                ;SAVE B AND C REGISTERS
; PUSH H                                ;SAVE HL REGISTERS
; .....
; perform disk read at this point, branch to
; label START if an error occurs
; .....
POP H                                  ;RECOVER HL
POP B                                  ;RECOVER B AND C REGISTERS
RET                                    ;BACK TO MAIN PROGRAM

END START

```

This program is assembled and listed in Appendix B for reference purposes, with an assumed origin of 100H. The hexadecimal operation codes that are listed on the left may be useful if the program has to be entered through the panel switches.

The PUTSYS program can be constructed from GETSYS by changing only a few operations in the GETSYS program given above, as shown in Appendix C. The register pair HL becomes the dump address (next address to write), and operations upon these registers do not change within the program. The READSEC subroutine is replaced by a WRITESEC subroutine, which performs the opposite function: data from address HL are written to the track given by register B and sector given by register C. It is often useful to combine GETSYS and PUTSYS into a single program during the test and development phase, as shown in Appendix C.

4A24H+b	JMP SETDMA	; SET DMA ADDRESS
4A27H+b	JMP READ	; READ SELECTED SECTOR
4A2AH+b	JMP WRITE	; WRITE SELECTED SECTOR
4A2DH+b	JMP LISTST	; RETURN LIST STATUS
4A30H+b	JMP SECTTRAN	; SECTOR TRANSLATE SUBROUTINE

Each jump address corresponds to a particular subroutine that performs the specific function, as outlined below. There are three major divisions in the jump table: the system (re)initialization, which results from calls on BOOT and WBOOT; simple character I/O performed by calls on CONST, CONIN, CONOUT, LIST, PUNCH, READER, and LISTST; and diskette I/O performed by calls on HOME, SELDSK, SETTRK, SETSEC, SETDMA, READ, WRITE, and SECTTRAN.

All simple character I/O operations are assumed to be performed in ASCII, upper and lower case, with high order (parity bit) set to zero. An end-of-file condition for an input device is given by an ASCII control-z (1AH). Peripheral devices are seen by CP/M as "logical" devices and are assigned to physical devices within the BIOS.

To operate, the BDOS needs only the CONST, CONIN, and CONOUT subroutines (LIST, PUNCH, and READER may be used by PIP, but not the BDOS). Further, the LISTST entry is currently used only by DESPOOL, the print spooling utility. Thus, the initial version of CBIOS may have empty subroutines for the remaining ASCII devices.

The characteristics of each device are

CONSOLE	The principal interactive console that communicates with the operator, accessed through CONST, CONIN, and CONOUT; typically, the CONSOLE is a device such as a CRT or teletype.
LIST	The principal listing device, if it exists on the user's system, is usually a hard-copy device, such as a printer or teletype.
PUNCH	The principal tape punching device, if it exists, is normally a high-speed paper tape punch or teletype.
READER	The principal tape reading device, such as a simple optical reader or teletype.

A single peripheral can be assigned as the LIST, PUNCH, and READER device simultaneously. If no peripheral device is assigned as the LIST, PUNCH, or READER device, the CBIOS created by the user may give an appropriate error message so that the system does not "hang" if the device is accessed by PIP or some other user program. Alternately, the PUNCH and LIST routines can just simply return, and the READER routine can return with a 1AH (ctl-Z) in register A to indicate immediate end-of-file.

For added flexibility, the user can optionally implement the "IOBYTE" function, which allows reassignment of physical and logical devices. The IOBYTE function creates a mapping of logical to physical devices that can be altered during CP/M processing (the user should see the STAT command). The definition of the IOBYTE function corresponds to the Intel standard as follows: a single location in memory (currently location 0003H) is maintained, called IOBYTE, which defines the logical to physical device mapping that is in effect at a particular time. The mapping is performed by splitting the

IOBYTE into four distinct fields of two bits each, called the CONSOLE, READER, PUNCH, and LIST fields, as shown below.

	most significant		least significant	
IOBYTE AT 003H	LIST	PUNCH	READER	CONSOLE
	bits 6, 7	bits 4, 5	bits 2, 3	bits 0, 1

The value in each field can be in the range 0-3, defining the assigned source or destination of each logical device. The values that can be assigned to each field are given below

CONSOLE field (bits 0,1)

- 0 console is assigned to the console printer device (TTY:)
- 1 console is assigned to the CRT device (CRT:)
- 2 batch mode: use the READER as the CONSOLE input, and the LIST device as the CONSOLE output (BAT:)
- 3 user defined console device (UC1:)

READER field (bits 2,3)

- 0 READER is the teletype device (TTY:)
- 1 READER is the high speed reader device (PTR:)
- 2 user defined reader # 1 (UR1:)
- 3 user defined reader # 2 (UR2:)

PUNCH field (bits 4,5)

- 0 PUNCH is the teletype device (TTY:)
- 1 PUNCH is the high speed punch device (PTP:)
- 2 user defined punch # 1 (UP1:)
- 3 user defined punch # 2 (UP2:)

LIST field (bits 6,7)

- 0 LIST is the teletype device (TTY:)
- 1 LIST is the CRT device (CRT:)
- 2 LIST is the line printer device (LPT:)
- 3 user defined list device (UL1:)

The implementation of the IOBYTE is optional and affects only the organization of the CBIOS. No CP/M systems use the IOBYTE (although they tolerate the existence of the IOBYTE at location 0003H), except for PIP, which allows access to the physical devices, and STAT, which allows logical-physical assignments to be made or displayed (for more information, the user should see Chapter 1). In any case the IOBYTE implementation should be omitted until the basic CBIOS is fully implemented and tested; then the user should add the IOBYTE to increase the facilities.

Disk I/O is always performed through a sequence of calls on the various disk access subroutines that set up the disk number to access, the track and sector on a particular disk, and the direct memory access (DMA) address involved in the I/O operation. After all these parameters have been set up, a call is made to the READ or WRITE function to perform the actual I/O operation. There is often a single call to SELDSK to select a disk drive, followed by a number of read or write operations to the selected disk before selecting another drive for subsequent operations. Similarly, there may be a single call to set the DMA address, followed by several calls that read or write from the selected DMA address before the DMA address is changed. The track and sector subroutines are always called before the READ or WRITE operations are performed.

The READ and WRITE routines should perform several retries (10 is standard) before reporting the error condition to the BDOS. If the error condition is returned to the BDOS, it will report the error to the user. The HOME subroutine may or may not

actually perform the track 00 seek, depending upon controller characteristics; the important point is that track 00 has been selected for the next operation and is often treated in exactly the same manner as SETTRK with a parameter of 00.

The exact responsibilities of each entry point subroutine are given below.

BOOT The BOOT entry point gets control from the cold start loader and is responsible for basic system initialization, including sending a sign-on message (which can be omitted in the first version). If the IOBYTE function is implemented, it must be set at this point. The various system parameters that are set by the WBOOT entry point must be initialized, and control is transferred to the CCP at 3400+b for further processing. Note that register C must be set to zero to select drive A.

WBOOT The WBOOT entry point gets control when a warm start occurs. A warm start is performed whenever a user program branches to location 0000H, or when the CPU is reset from the front panel. The CP/M system must be loaded from the first two tracks of drive A up to, but not including, the BIOS (or CBIOS, if the user has completed the patch). System parameters must be initialized as shown below:

location 0,1,2	Set to JMP WBOOT for warm starts (000H: JMP 4A03H+b)
location 3	Set initial value of IOBYTE, if implemented in the CBIOS
location 4	High nibble = current user no; low nibble = current drive
location 5,6,7	Set to JMP BDOS, which is the primary entry point to CP/M for transient programs. (0005H: JMP 3C06H+b)

(The user should refer to Section 6.9 for complete details of page zero use.) Upon completion of the initialization, the WBOOT program must branch to the CCP at 3400H+b to (re)start the system. Upon entry to the CCP, register C is set to the drive to select after system initialization. The WBOOT routine should read location 4 in memory, verify that it is a legal drive, and pass it to the CCP in register C.

CONST The user should sample the status of the currently assigned console device and return 0FFH in register A if a character is ready to read and 00H in register A if no console characters are ready.

CONIN The next console character is read into register A, and the parity bit is set (high order bit) to zero. If no console character is ready, the user waits until a character is typed before returning.

CONOUT	The user sends the character from register C to the console output device. The character is in ASCII, with high order parity bit set to zero. The user may want to include a time-out on a line feed or carriage return, if the console device requires some time interval at the end of the line (such as a TI Silent 700 terminal). The user can filter out control characters that cause the console device to react in a strange way (a control-z causes the Lear Seigler terminal to clear the screen, for example).
LIST	The user sends the character from register C to the currently assigned listing device. The character is in ASCII with zero parity bit.
PUNCH	The user sends the character from register C to the currently assigned punch device. The character is in ASCII with zero parity.
READER	The user reads the next character from the currently assigned reader device into register A with zero parity (high order bit must be zero); an end-of-file condition is reported by returning an ASCII control-z(1AH).
HOME	The user moves the disk head of the currently selected disk (initially disk A) to the track 00 position. If the controller allows access to the track 0 flag from the drive, the head is stepped until the track 0 flag is detected. If the controller does not support this feature, the HOME call is translated into a call to SETTRK with a parameter of 0.
SELDSK	The user selects the disk drive given by register C for further operations, where register C contains 0 for drive A, 1 for drive B, and so on up to 15 for drive P (the standard CP/M distribution version supports four drives). On each disk select, SELDSK must return in HL the base address of a 16-byte area, called the Disk Parameter Header, described in Section 6.10. For standard floppy disk drives, the contents of the header and associated tables do not change; thus, the program segment included in the sample CBIOS performs this operation automatically. If there is an attempt to select a nonexistent drive, SELDSK returns HL=0000H as an error indicator. Although SELDSK must return the header address on each call, it is advisable to postpone the physical disk select operation until an I/O function (seek, read, or write) is actually performed, since disk selects often occur without ultimately performing any disk I/O, and many controllers will unload the head of the current disk before selecting the new drive. This would cause an excessive amount of noise and disk wear. The least significant bit of register E is zero if this is the first occurrence of the drive select since the last cold or warm start.
SETTRK	Register BC contains the track number for subsequent disk accesses on the currently selected drive. The sector number in BC is the same as the number returned from the SECTTRAN entry point. The user can choose to seek the selected track at

this time or delay the seek until the next read or write actually occurs. Register BC can take on values in the range 0-76 corresponding to valid track numbers for standard floppy disk drives and 0-65535 for nonstandard disk subsystems.

SETSEC Register BC contains the sector number (1 through 26) for subsequent disk accesses on the currently selected drive. The sector number in BC is the same as the number returned from the SECTTRAN entry point. The user can choose to send this information to the controller at this point or delay sector selection until a read or write operation occurs.

SETDMA Register BC contains the DMA (disk memory access) address for subsequent read or write operations. For example, if B = 00H and C = 80H when SETDMA is called, all subsequent read operations read their data into 80H through 0FFH and all subsequent write operations get their data from 80H through 0FFH, until the next call to SETDMA occurs. The initial DMA address is assumed to be 80H. The controller need not actually support direct memory access. If, for example, all data transfers are through I/O ports, the CBIOS that is constructed will use the 128-byte area starting at the selected DMA address for the memory buffer during the subsequent read or write operations.

READ Assuming the drive has been selected, the track has been set, the sector has been set, and the DMA address has been specified, the READ subroutine attempts to read one sector based upon these parameters and returns the following error codes in register A:

- 0 no errors occurred
- 1 nonrecoverable error condition occurred

Currently, CP/M responds only to a zero or nonzero value as the return code. That is, if the value in register A is 0, CP/M assumes that the disk operation was completed properly. If an error occurs, however, the CBIOS should attempt at least 10 retries to see if the error is recoverable. When an error is reported the BDOS will print the message "BDOS ERR ON x: BAD SECTOR". The operator then has the option of typing carriage-return to ignore the error, or ctl-C to abort.

WRITE The user writes the data from the currently selected DMA address to the currently selected drive, track, and sector. For floppy disks, the data should be marked as "nondeleted data" to maintain compatibility with other CP/M systems. The error codes given in the READ command are returned in register A, with error recovery attempts as described above.

LISTST The user returns the ready status of the list device used by the DESPOOL program to improve console response during its operation. The value 00 is returned in A if the list device is not ready to accept a character and 0FFH if a character can be sent

to the printer. A 00 value should be returned if LIST status is not implemented.

SECTRAN

The user performs logical to physical sector translation to improve the overall response of CP/M. Standard CP/M systems are shipped with a "skew factor" of 6, where six physical sectors are skipped between each logical read operation. This skew factor allows enough time between sectors for most programs to load their buffers without missing the next sector. In particular computer systems that use fast processors, memory, and disk subsystems, the skew factor may be changed to improve overall response. However, the user should maintain a single density IBM-compatible version of CP/M for information transfer into and out of the computer system, using a skew factor of 6. In general, SECTRAN receives a logical sector number relative to zero in BC and a translate table address in DE. The sector number is used as an index into the translate table, with the resulting physical sector number in HL. For standard systems, the table and indexing code is provided in the CBIOS and need not be changed.

6.7 A Sample BIOS

The program shown in Appendix B can serve as a basis for a user's first BIOS. The simplest functions are assumed in this BIOS, so that the user can enter it through a front panel, if absolutely necessary. The user must alter and insert code into the subroutines for CONST, CONIN, CONOUT, READ, WRITE, and WAITIO subroutines. Storage is reserved for user-supplied code in these regions. The scratch area reserved in page zero (see section 6.9) for the BIOS is used in this program, so that it could be implemented in ROM, if desired.

Once operational, this skeletal version can be enhanced to print the initial sign-on message and perform better error recovery. The subroutines for LIST, PUNCH, and READER can be filled out and the IOBYTE function can be implemented.

6.8 A Sample Cold Start Loader

The program shown in Appendix E can serve as a basis for a cold start loader. The disk read function must be supplied by the user, and the program must be loaded somehow starting at location 0000. Space is reserved for the patch code so that the total amount of storage required for the cold start loader is 128 bytes. Eventually, the user will probably want to get this loader onto the first disk sector (track 0, sector 1) and cause the controller to load it into memory automatically upon system start up. Alternatively, the cold start loader can be placed into ROM, and above the CP/M system. In this case, it will be necessary to originate the program at a higher address and key in a jump instruction at system start up that branches to the loader. Subsequent warm starts will not require this key-in operation, since the entry point WBOOT gets control thus bringing the system in from disk automatically. The skeletal cold start loader has minimal error recover, which may be enhanced in later versions.

6.9 Reserved Locations in Page Zero

Main memory page zero, between locations 00H and 0FFH, contains several segments of code and data that are used during CP/M processing. The code and data areas are given below for reference

Locations from to	Contents
0000H-0002H	Contains a jump instruction to the warm start entry point at location 4A03H+b. This allows a simple programmed restart (JMP 0000H) or manual restart from the front panel.
0003H-0003H	Contains the Intel standard IOBYTE, which is optionally included in the user's CBIOS, as described in Section 6.6.
0004H-0004H	Current default drive number (0=A,...,15=P).
0005H-0007H	Contains a jump instruction to the BDOS and serves two purposes: JMP 0005H provides the primary entry point to the BDOS, as described in Chapter 5, and LHL 0006H brings the address field of the instruction to the HL register pair. This value is the lowest address in memory used by CP/M (assuming the CCP is being overlaid). The DDT program will change the address field to reflect the reduced memory size in debug mode.
0008H-0027H	(Interrupt locations 1 through 5 not used.)
0030H-0037H	(Interrupt location 6, not currently used; reserved.)
0038H-003AH	Restart 7; contains a jump instruction into the DDT or SID program when running in debug mode for programmed breakpoints, but is not otherwise used by CP/M.
003BH-003FH	(Not currently used; reserved.)
0040H-004FH	A 16-byte area reserved for scratch by CBIOS, but is not used for any purpose in the distribution version of CP/M.
0050H-005BH	(Not currently used; reserved.)
005CH-007CH	Default file control block produced for a transient program by the Console Command Processor.
007DH-007FH	Optional default random record position.
0080H-00FFH	Default 128-byte disk buffer (also filled with the command line when a transient is loaded under the CCP).

This information is set up for normal operation under the CP/M system, but can be overwritten by a transient program if the BDOS facilities are not required by the transient.

If, for example, a particular program performs only simple I/O and must begin execution at location 0, it can first be loaded into the TPA, using normal CP/M facilities, with a small memory move program that gets control when loaded (the memory move program must get control from location 0100H, which is the assumed beginning of all transient programs). The move program can then proceed to move the entire memory image down to location 0 and pass control to the starting address of the memory load. If the BIOS is overwritten or if location 0 (containing the warm start entry point) is overwritten, the operator must bring the CP/M system back into memory with a cold start sequence.

6.10 Disk Parameter Tables

Tables are included in the BIOS that describe the particular characteristics of the disk subsystem used with CP/M. These tables can be either hand-coded, as shown in the sample CBIOS in Appendix B, or automatically generated using the DISKDEF macro library, as shown in Appendix F. The purpose here is to describe the elements of these tables.

In general, each disk drive has an associated (16-byte) disk parameter header that contains information about the disk drive and provides a scratchpad area for certain BDOS operations. The format of the disk parameter header for each drive is shown below.

Disk Parameter Header							
XLT	0000	0000	0000	DIRBUF	DPB	CSV	ALV
16b	16b	16b	16b	16b	16b	16b	16b

where each element is a word (16-bit) value. The meaning of each Disk Parameter Header (DPH) element is

XLT	Address of the logical to physical translation vector, if used for this particular drive, or the value 0000H if no sector translation takes place (i.e., the physical and logical sector numbers are the same). Disk drives with identical sector skew factors share the same translate tables.
0000	Scratchpad values for use within the BDOS (initial value is unimportant).
DIRBUF	Address of a 128-byte scratchpad area for directory operations within BDOS. All DPHs address the same scratchpad area.
DPB	Address of a disk parameter block for this drive. Drives with identical disk characteristics address the same disk parameter block.
CSV	Address of a scratchpad area used for software check for changed disks. This address is different for each DPH.
ALV	Address of a scratchpad area used by the BDOS to keep disk storage allocation information. This address is different for each DPH.

Given n disk drives, the DPHs are arranged in a table whose first row of 16 bytes corresponds to drive 0, with the last row corresponding to drive n-1. The table thus appears as

DPBASE:

00	XLT 00	0000	0000	0000	DIRBUF	DBP 00	CSV 00	ALV 00
01	XLT 01	0000	0000	0000	DIRBUF	DBP 01	CSV 01	ALV 01
(and so on through)								
n-1	XLTn-1	0000	0000	0000	DIRBUF	DBPn-1	CSVn-1	ALVn-1

where the label DPBASE defines the base address of the DPH table.

A responsibility of the SELDSK subroutine is to return the base address of the DPH for the selected drive. The following sequence of operations returns the table address, with a 0000H returned if the selected drive does not exist.

```

NDISKS      EQU      4      ;NUMBER OF DISK DRIVES
.....
SELDISK:    ;SELECT DISK GIVEN BY BC
            LXI      H,0000H ;ERROR CODE
            MOV      A,C      ;DRIVE OK?
            CPI      NDISKS  ;CY IF SO
            RNC      ;RET IF ERROR
            ;NO ERROR, CONTINUE
            MOV      L,C      ;LOW(DISK)
            MOV      H,B      ;HIGH(DISK)
            DAD      H        ;*2
            DAD      H        ;*4
            DAD      H        ;*8
            DAD      H        ;*16
            LXI      D,DPBASE;FIRST DPH
            DAD      D        ;DPH(DISK)
            RET

```

The translation vectors (XLT 00 through XLTn-1) are located elsewhere in the BIOS, and simply correspond one-for-one with the logical sector numbers zero through the sector count 1. The Disk Parameter Block (DPB) for each drive is more complex. A particular DPB, which is addressed by one or more DPHs, takes the general form

SPT	BSH	BLM	EXM	DSM	DRM	AL0	AL1	CKS	OFF
16b	8b	8b	8b	16b	16b	8b	8b	16b	16b

where each is a byte or word value, as shown by the 8b or 16b indicator below the field.

SPT is the total number of sectors per track.

BSH is the data allocation block shift factor, determined by the data block allocation size.

BLM is the data allocation block mask ($2^{[BSH-1]}$).

EXM is the extent mask, determined by the data block allocation size and the number of disk blocks.

DSM determines the total storage capacity of the disk drive.

DRM	determines the total number of directory entries that can be stored on this drive. (AL0,AL1 determine reserved directory blocks.)
CKS	is the size of the directory check vector.
OFF	is the number of reserved tracks at the beginning of the (logical) disk.

The values of BSH and BLM determine (implicitly) the data allocation size BLS, which is not an entry in the DPB. Given that the designer has selected a value for BLS, the values of BSH and BLM are shown in the tabulation below.

BLS	BSH	BLM
1024	3	7
2048	4	15
4096	5	31
8192	6	63
16384	7	127

where all values are in decimal. The value of EXM depends upon both the BLS and whether the DSM value is less than 256 or greater than 255. For $DSM < 256$ the value of EXM is given by:

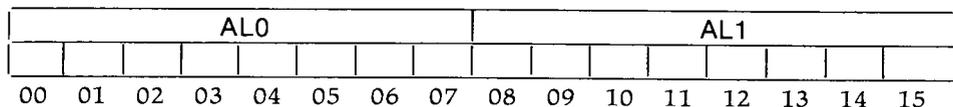
BLS	EXM
1024	0
2048	1
4096	3
8192	7
16384	15

For $DSM > 255$ the value of EXM is given by:

BLS	EXM
1024	N/A
2048	0
4096	1
8192	3
16384	7

The value of DSM is the maximum data block number supported by this particular drive, measured in BLS units. The product BLS times (DSM+1) is the total number of bytes held by the drive and, of course, must be within the capacity of the physical disk, not counting the reserved operating system tracks.

The DRM entry is the one less than the total number of directory entries that can take on a 16-bit value. The values of AL0 and AL1, however, are determined by DRM. The values AL0 and AL1 can together be considered a string of 16-bits, as shown below.



where position 00 corresponds to the high order bit of the byte labeled AL0 and 15 corresponds to the low order bit of the byte labeled AL1. Each bit position reserves a data block for number of directory entries, thus allowing a total of 16 data blocks to be

assigned for directory entries (bits are assigned starting at 00 and filled to the right until position 15). Each directory entry occupies 32 bytes, resulting in the tabulation below.

BLS	Directory Entries
1024	32 times # bits
2048	64 times # bits
4096	128 times # bits
8192	256 times # bits
16384	512 times # bits

Thus, if DRM = 127 (128 directory entries) and BLS = 1024, there are 32 directory entries per block, requiring 4 reserved blocks. In this case, the 4 high order bits of AL0 are set, resulting in the values AL0 = 0F0H and AL1 = 00H.

The CKS value is determined as follows: if the disk drive media is removable, then $CKS = (DRM+1)/4$, where DRM is the last directory entry number. If the media are fixed, then set CKS = 0 (no directory records are checked in this case).

Finally, the OFF field determines the number of tracks that are skipped at the beginning of the physical disk. This value is automatically added whenever SETTRK is called and can be used as a mechanism for skipping reserved operating system tracks or for partitioning a large disk into smaller segmented sections.

To complete the discussion of the DPB, several DPHs can address the same DPB if their drive characteristics are identical. Further, the DPB can be dynamically changed when a new drive is addressed by simply changing the pointer in the DPH since the BDOS copies the DPB values to a local area whenever the SELDSK function is invoked.

Returning back to the DPH for a particular drive, the two address values CSV and ALV remain. Both addresses reference an area of uninitialized memory following the BIOS. The areas must be unique for each drive, and the size of each area is determined by the values in the DPB.

The size of the area addressed by CSV is CKS bytes, which is sufficient to hold the directory check information for this particular drive. If $CKS = (DRM+1)/4$, one must reserve $(DRM+1)/4$ bytes for directory check use. If CKS = 0, no storage is reserved.

The size of the area addressed by ALV is determined by the maximum number of data blocks allowed for this particular disk and is computed as $(DSM/8)+1$.

The CBIOS shown in Appendix B demonstrates an instance of these tables for standard 8-inch single density drives. It may be useful to examine this program and compare the tabular values with the definitions given above.

6.11 The DISKDEF Macro Library

A macro library is shown in Appendix F, called DISKDEF, which greatly simplifies the table construction process. One must have access to the MAC macro assembler, of course, to use the DISKDEF facility, while the macro library is included with all CP/M 2 distribution disks.

A BIOS disk definition consists of the following sequence of macro statements:

MACLIB	DISKDEF
.....	
DISKS	n
DISKDEF	0,...
DISKDEF	1,...

```
.....  
DISKDEF          n-1  
.....  
ENDEF
```

where the MACLIB statement loads the DISKDEF.LIB file (on the same disk as the BIOS) into MAC's internal tables. The DISKS macro call follows, which specifies the number of drives to be configured with the user's system, where n is an integer in the range 1 to 16. A series of DISKDEF macro calls then follow that define the characteristics of each logical disk, 0 through n-1 (corresponding to logical drives A through P). The DISKS and DISKDEF macros generate the in-line fixed data tables described in the previous section and thus must be placed in a nonexecutable portion of the BIOS, typically directly following the BIOS jump vector.

The remaining portion of the BIOS is defined following the DISKDEF macros, with the ENDEF macro call immediately preceding the END statement. The ENDEF (End of Diskdef) macro generates the necessary uninitialized RAM areas, which are located in memory above the BIOS.

The form of the DISKDEF macro call is

```
DISKDEF dn,fsc,lsc,[skf],bls,dks,dir,cks,ofs,[0]
```

where

dn	is the logical disk number, 0 to n-1.
fsc	is the first physical sector number (0 or 1).
lsc	is the last sector number.
skf	is the optional sector skew factor.
bls	is the data allocation block size.
dks	is the number of blocks on the disk.
dir	is the number of directory entries.
cks	is the number of "checked" directory entries.
ofs	is the track offset to logical track 00.
[0]	is an optional 1.4 compatibility flag.

The value dn is the drive number being defined with this DISKDEF macro invocation. The fsc parameter accounts for differing sector numbering systems and is usually 0 or 1. The lsc is the last numbered sector on a track. When present, the skf parameter defines the sector skew factor, which is used to create a sector translation table according to the skew.

If the number of sectors is less than 256, a single-byte table is created, otherwise each translation table element occupies two bytes. No translation table is created if the skf parameter is omitted (or equal to 0). The bls parameter specifies the number of bytes allocated to each data block, and takes on the values 1024, 2048, 4096, 8192, or 16384. Generally, performance increases with larger data block sizes since there are fewer directory references and logically connected data records are physically close on the disk. Further, each directory entry addresses more data and the BIOS-resident ram space is reduced.

The dks parameter specifies the total disk size in bls units. That is, if the bls = 2048 and dks = 1000, the total disk capacity is 2,048,000 bytes. If dks is greater than 255, the block size parameter bls must be greater than 1024. The value of dir is the total number of directory entries, which may exceed 255, if desired. The cks parameter determines the

number of directory items to check on each directory scan and is used internally to detect changed disks during system operation, where an intervening cold or warm start has not occurred (when this situation is detected, CP/M automatically marks the disk read/only so that data are not subsequently destroyed).

As stated in the previous section, the value of `cks = dir` when the medium is easily changed, as is the case with a floppy disk subsystem. If the disk is permanently mounted, the value of `cks` is typically 0, since the probability of changing disks without a restart is low. The `ofs` value determines the number of tracks to skip when this particular drive is addressed, which can be used to reserve additional operating system space or to simulate several logical drives on a single large capacity physical drive. Finally, the `[0]` parameter is included when file compatibility is required with versions of 1.4 that have been modified for higher density disks. This parameter ensures that only 16K is allocated for each directory record, as was the case for previous versions. Normally, this parameter is not included.

For convenience and economy of table space, the special form

```
DISKDEF      i,j
```

gives disk `i` the same characteristics as a previously defined drive `j`. A standard four-drive single density system, which is compatible with version 1.4, is defined using the following macro invocations:

```
DISKS          4
DISKDEF        0,1,26,6,1024,243,64,64,2
DISKDEF        1,0
DISKDEF        2,0
DISKDEF        3,0
...
ENDEF
```

with all disks having the same parameter values of 26 sectors per track (numbered 1 through 26), with 6 sectors skipped between each access, 1024 bytes per data block, 243 data blocks for a total of 243K-byte disk capacity, 64 checked directory entries, and two operating system tracks.

The `DISKS` macro generates `n` DPHs, starting at the DPH table address `DPBASE` generated by the macro. Each disk header block contains sixteen bytes, as described above, and correspond one-for-one to each of the defined drives. In the four-drive standard system, for example, the `DISKS` macro generates a table of the form:

```
DPBASE        EQU $
DPE0:         DW  XLT0,0000H,0000H,0000H,DIRBUF,DPB0,CSV0,ALV0
DPE1:         DW  XLT0,0000H,0000H,0000H,DIRBUF,DPB0,CSV1,ALV1
DPE2:         DW  XLT0,0000H,0000H,0000H,DIRBUF,DPB0,CSV2,ALV2
DPE3:         DW  XLT0,0000H,0000H,0000H,DIRBUF,DPB0,CSV3,ALV3
```

where the DPH labels are included for reference purposes to show the beginning table addresses for each drive 0 through 3. The values contained within the DPH are described in detail in the previous section. The check and allocation vector addresses are generated by the `ENDEF` macro in the ram area following the BIOS code and tables.

The user should note that if the `skf` (skew factor) parameter is omitted (or equal to 0), the translation table is omitted and a 0000H value is inserted in the XLT position of the DPH for the disk. In a subsequent call to perform the logical to physical translation, `SECTTRAN` receives a translation table address of `DE = 0000H` and simply returns the original logical sector from `BC` in the `HL` register pair. A translate table is constructed when the `skf` parameter is present, and the (nonzero) table address is placed into the

corresponding DPHs. The tabulation shown below, for example, is constructed when the standard skew factor $skf = 6$ is specified in the DISKDEF macro call:

```

XLT0:      DB      1,7,13,19,25,5,11,17,23,3,9,15,21
           DB      2,8,14,20,26,6,12,18,24,4,10,16,22

```

Following the ENDEF macro call, a number of uninitialized data areas are defined. These data areas need not be a part of the BIOS that is loaded upon cold start, but must be available between the BIOS and the end of memory. The size of the uninitialized RAM area is determined by EQU statements generated by the ENDEF macro. For a standard four-drive system, the ENDEF macro might produce

```

4C72 =      BEGDAT EQU $
           (data areas)
4DB0 =      ENDDAT EQU $
013C =      DATSIZ EQU $-BEGDAT

```

which indicates that uninitialized RAM begins at location 4C72H, ends at 4DB0H-1, and occupies 013CH bytes. The user must ensure that these addresses are free for use after the system is loaded.

After modification, the user can utilize the STAT program to check drive characteristics, since STAT uses the disk parameter block to decode the drive information. The STAT command form

```
STAT d:DSK:
```

decodes the disk parameter block for drive d (d=A,...,P) and displays the values shown below.

```

r: 128-byte record capacity
k: kilobyte drive capacity
d: 32-byte directory entries
c: checked directory entries
e: records/extent
b: records/block
s: sectors/track
t: reserved tracks

```

Three examples of DISKDEF macro invocations are shown below with corresponding STAT parameter values (the last produces a full 8-megabyte system).

```

r=4096,      DISKDEF 0,1,58,,2048,256,128,128,2
           k=512, d=128, c=128, e=256, b=16, s=58, t=2
r=16384,     DISKDEF 0,1,58,,2048,1024,300,0,2
           k=2048, d=300, c=0, e=128, b=16, s=58, t=2
r=65536,     DISKDEF 0,1,58,,16384,512,128,128,2
           k=8192, d=128, c=128, e=1024, b=128, s=58, t=2

```

6.12 Sector Blocking and Deblocking

Upon each call to the BIOS WRITE entry point, the CP/M BDOS includes information that allows effective sector blocking and deblocking where the host disk subsystem has a sector size that is a multiple of the basic 128-byte unit. The purpose here is to present a general-purpose algorithm that can be included within the BIOS and that uses the BDOS information to perform the operations automatically.

On each call to WRITE, the BDOS provides the following information in register C:

0	=	normal sector write
1	=	write to directory sector
2	=	write to the first sector of a new data block

Condition 0 occurs whenever the next write operation is into a previously written area, such as a random mode record update, when the write is to other than the first sector of an unallocated block, or when the write is not into the directory area. Condition 1 occurs when a write into the directory area is performed. Condition 2 occurs when the first record (only) of a newly allocated data block is written. In most cases, application programs read or write multiple 128-byte sectors in sequence; thus, there is little overhead involved in either operation when blocking and deblocking records, since preread operations can be avoided when writing records.

Appendix G lists the blocking and deblocking algorithms in skeletal form (this file is included on your CP/M disk). Generally, the algorithms map all CP/M sector read operations onto the host disk through an intermediate buffer that is the size of the host disk sector. Throughout the program, values and variables that relate to the CP/M sector involved in a seek operation are prefixed by sek, while those related to the host disk system are prefixed by hst. The equate statements beginning on line 29 of Appendix G define the mapping between CP/M and the host system, and must be changed if other than the sample host system is involved.

The entry points BOOT and WBOOT must contain the initialization code starting on line 57, while the SELDSK entry point must be augmented by the code starting on line 65. The user should note that although the SELDSK entry point computes and returns the Disk Parameter Header address, it does not physically select the host disk at this point (it is selected later at READHST or WRITEHST). Further, SETTRK, SETTRK, and SETDMA simply store the values, but do not take any other action at this point. SECTRAN performs a trivial function of returning the physical sector number.

The principal entry points are READ and WRITE, starting on lines 110 and 125, respectively. These subroutines take the place of your previous READ and WRITE operations.

The actual physical read or write takes place at either WRITEHST or READHST, where all values have been prepared: hstdsk is the host disk number, hstrk is the host track number, and hstsec is the host sector number (which may require translation to a physical sector number). The user must insert code at this point that performs the full host sector read or write into or out of the buffer at hstbuf of length hstsiz. All other mapping functions are performed by the algorithms.

This particular algorithm was tested using an 80-megabyte hard disk unit that was originally configured for 128-byte sectors, producing approximately 35 megabytes of formatted storage. When configured for 512-byte host sectors, usable storage increased to 57 megabytes, with a corresponding 400% improvement in overall response. In this situation, there is no apparent overhead involved in deblocking sectors, with the advantage that user programs still maintain 128-byte sectors. This is primarily because of the information provided by the BDOS, which eliminates the necessity for preread operations.

CP/M[®]
OPERATING SYSTEM
MANUAL

Appendices A-G

 **DIGITAL RESEARCH[™]**
P.O. Box 579
Pacific Grove, California 93950

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Appendix A: The MDS Basic I/O System (BIOS)

```

1           ;           mds-800 i/o drivers for cp/m 2.2
2           ;           (four drive single density version)
3           ;
4           ;           version 2.2 february, 1980
5           ;
6   0016 =   vers       equ    22           ;version 2.2
7           ;
8           ;           copyright (c) 1980
9           ;           digital research
10          ;           box 579, pacific grove
11          ;           california, 93950
12          ;
13          ;
14   ffff =   true       equ    offfh       ;value of "true"
15   0000 =   false      equ    not true    ;"false"
16   0000 =   test       equ    false      ; true if test bios
17          ;
18          ;           if          test
19          bias       equ    03400h       ;base of ccp in test system
20          ;           endif
21          ;           if          not test
22   0000 =   bias       equ    0000h       ;generate relocatable cp/m system
23          ;           endif
24          ;
25   1600 =   patch      equ    1600h
26          ;
27   1600          org    patch
28   0000 =   cpmb       equ    $-patch    ;base of cpm console processor
29   0806 =   bdos       equ    806h+cpmb ;basic dos (resident portion)

```

```

30      1600 =          cpml    equ      $-cpmb    ;length (in bytes) of cpm system
31      002c =          nsects  equ      cpml/128   ;number of sectors to load
32      0002 =          offset  equ      2         ;number of disk tracks used by cp/m
33      0004 =          cdisk   equ      0004h     ;address of last logged disk on warm start
34      0080 =          buff    equ      0080h     ;default buffer address
35      000a =          retry   equ      10        ;max retries on disk i/o before error
36      ;
37      ;          perform following functions
38      ;          boot      cold start
39      ;          wboot    warm start (save i/o byte)
40      ;          (boot and wboot are the same for mds)
41      ;          const   console status
42      ;          ;          reg-a = 00 if no character ready
43      ;          ;          reg-a = ff if character ready
44      ;          conin   console character in (result in reg-a)
45      ;          conout  console character out (char in reg-c)
46      ;          list   list out (char in reg-c)
47      ;          punch  punch out (char in reg-c)
48      ;          reader  paper tape reader in (result to reg-a)
49      ;          home   move to track 00
50      ;
51      ;          (the following calls set-up the io parameter block for the
52      ;          mds, which is used to perform subsequent reads and writes)
53      ;          seldsk  select disk given by reg-c (0, 1, 2. . .)
54      ;          settrk  set track address (0, . . . 76) for subsequent read/write
55      ;          setsec  set sector address (1, . . . , 26) for subsequent read/write
56      ;          setdma  set subsequent dma address (initially 80h)
57      ;
58      ;          (read and write assume previous calls to set up the io parameters)
59      ;          read   read track/sector to preset dma address
60      ;          write  track/sector from preset dma address
61      ;
62      ;          jump vector for individual routines
63      1600 c3b316     jmp      boot

```

```

64      1603 c3c316      wboote: jmp      wboot
65      1606 c36117      jmp      const
66      1609 c36417      jmp      conin
67      160c c36a17      jmp      conout
68      160f c36d17      jmp      list
69      1612 c37217      jmp      punch
70      1615 c37517      jmp      reader
71      1618 c37817      jmp      home
72      161b c37d17      jmp      seldsk
73      161e c3a717      jmp      settrk
74      1621 c3ac17      jmp      setsec
75      1624 c3bb17      jmp      setdma
76      1627 c3c117      jmp      read
77      162a c3ca17      jmp      write
78      162d c37017      jmp      listst      ;list status
79      1630 c3b117      jmp      sectran
80
81      ;
82      maclib diskdef      ;load the disk definition library
83      disks      4      ;four disks
84      1633+=      dpbase equ      $      ;base of disk parameter blocks
85      1633+82160000 dpe0: dw      xlt0, 0000h      ;translate table
86      1637+00000000 dw      0000h, 0000h      ;scratch area
87      163b+6e187316 dw      dirbuf, dpb0      ;dir buff, parm block
88      163f+0d19ee18 dw      csv0, alv0      ;check, alloc vectors
89      1643+82160000 dpe1: dw      xlt1, 0000h      ;translate table
90      1647+00000000 dw      0000h, 0000h      ;scratch area
91      164b+6e187316 dw      dirbuf, dpb1      ;dir buff, parm block
92      164f+3c191d19 dw      csv1, alv1      ;check, alloc vectors
93      1653+82160000 dpe2: dw      xlt2, 0000h      ;translate table
94      1657+00000000 dw      0000h, 0000h      ;scratch area
95      165b+6e187316 dw      dirbuf, dpb2      ;dir buff, parm block
96      165f+6b194c19 dw      csv2, alv2      ;check, alloc vectors
97      1663+82160000 dpe3: dw      xlt3, 0000h      ;translate table
98      1667+00000000 dw      0000h, 0000h      ;scratch area

```

98	166b+6e187316		dw	dirbuf, dpb3	;check, alloc block
99	166f+9a197b19		dw	csv3, alv3	;dir buff, parm vectors
100			diskdef	0, 1, 26, 6, 1024, 243, 64, 64,	offset
101	1673+=	dpb0	equ	\$;disk parm block
102	1673+1a00		dw	26	;sec per track
103	1675+03		db	3	;block shift
104	1676+07		db	7	;block mask
105	1677+00		db	0	;extnt mask
106	1678+f200		dw	242	;disk size-1
107	167a+3f00		dw	63	;directory max
108	167c+c0		db	192	;alloc0
109	167d+00		db	0	;alloc1
110	167e+1000		dw	16	;check size
111	1680+0200		dw	2	;offset
112	1682+=	xlt0	equ	\$;translate table
113	1682+01		db	1	
114	1683+07		db	7	
115	1684+0d		db	13	
116	1685+13		db	19	
117	1686+19		db	25	
118	1687+05		db	5	
119	1688+0b		db	11	
120	1689+11		db	17	
121	168a+17		db	23	
122	168b+03		db	3	
123	168c+09		db	9	
124	168d+0f		db	15	
125	168e+15		db	21	
126	168f+02		db	2	
127	1690+08		db	8	
128	1691+0e		db	14	
129	1692+14		db	20	
130	1693+1a		db	26	
131	1694+06		db	6	

```

132      1695+0c      db      12
133      1696+12      db      18
134      1697+18      db      24
135      1698+04      db      4
136      1699+0a      db      10
137      169a+10      db      16
138      169b+16      db      22
139      diskdef      1, 0
140      1673+=      dpb1      equ      dpb0      ;equivalent parameters
141      001f+=      als1      equ      als0      ;same allocation vector size
142      0010+=      css1      equ      css0      ;same checksum vector size
143      1682+=      xlt1      equ      xlt0      ;same translate table
144      diskdef      2, 0
145      1673+=      dpb2      equ      dpb0      ;equivalent parameters
146      001f+=      als2      equ      als0      ;same allocation vector size
147      0010+=      css2      equ      css0      ;same checksum vector size
148      1682+=      xlt2      equ      xlt0      ;same translate table
149      diskdef      3, 0
150      1673+=      dpb3      equ      dpb0      ;equivalent parameters
151      001f+=      als3      equ      als0      ;same allocation vector size
152      0010+=      css3      equ      css0      ;same checksum vector size
153      1682+=      xlt3      equ      xlt0      ;same translate table
154      ;      endif occurs at end of assembly
155      ;
156      ;      end of controller—independent code, the remaining subroutines
157      ;      are tailored to the particular operating environment, and must
158      ;      be altered for any system which differs from the intel mds.
159      ;
160      ;      the following code assumes the mds monitor exists at 0f800h
161      ;      and uses the i/o subroutines within the monitor
162      ;
163      ;      we also assume the mds system has four disk drives
164      00fd =      revrt      equ      0fdh      ;interrupt revert port
165      00fc =      intc       equ      0fch      ;interrupt mask port

```

```

166 00f3 = icon equ 0f3h ;interrupt control port
167 007E = inte equ 0111$1110b ;enable rst 0 (warm boot), rst 7 (monitor)
168 ;
169 ; mds monitor equates
170 f800 = mon80 equ 0f800h ;mds monitor
171 ff0f = rmon80 equ 0ff0fh ;restart mon80 (boot error)
172 f803 = ci equ 0f803h ;console character to reg-a
173 f806 = ri equ 0f806h ;reader in to reg-a
174 f809 = co equ 0f809h ;console char from c to console out
175 f80c = po equ 0f80ch ;punch char from c to punch device
176 f80f = lo equ 0f80fh ;list from c to list device
177 f812 = csts equ 0f812h ;console status 00/ff to register a
178 ;
179 ; disk ports and commands
180 0078 = base equ 78h ;base of disk command io ports
181 0078 = dstat equ base ;disk status (input)
182 0079 = rtype equ base+1 ;result type (input)
183 007b = rbyte equ base+3 ;result byte (input)
184 ;
185 0079 = ilow equ base+1 ;iopb low address (output)
186 007a = ihigh equ base+2 ;iopb high address (output)
187 ;
188 0004 = readf equ 4h ;read function
189 0006 = writf equ 6h ;write function
190 0003 = recal equ 3h ;recalibrate drive
191 0004 = iordy equ 4h ;i/o finished mask
192 000d = cr equ 0dh ;carriage return
193 000a = lf equ 0ah ;line feed
194 ;
195 signon: ;signon message: xxk cp/m vers y.y
196 169c 0d0a0a db cr, lf, lf
197 if test
198 db '32' ;32k example bios
199 endif

```

```

200          if      not test
201          169f 3030      db      '00'          ;memory size filled by relocater
202          endif
203          16a1 6b2043502f db      'k cp/m vers '
204          16ad 322e32      db      vers/10+'0', ', ' vers mod 10+'0'
205          16b0 0d0a00      db      cr, lf, 0
206          ;
207          boot:      ;print signon message and go to ccp
208          ;          (note: mds boot initialized iobyte at 0003h)
209          16b3 310001      lxi      sp, buff+80h
210          16b6 219c16      lxi      h, signon
211          16b9 cdd317      call     prmsg      ;print message
212          16bc af          xra      a          ;clear accumulator
213          16bd 320400      sta      cdisk     ;set initially to disk a
214          16c0 c30f17      jmp      gocpm      ;go to cp/m
215          ;
216          ;
217          wboot:; loader on track 0, sector 1, which will be skipped for warm
218          ;          read cp/m from disk—assuming there is a 128 byte cold start
219          ;          start
220          ;
221          16c3 318000      lxi      sp, buff      ;using dma—thus 80 thru ff available for stack
222          ;
223          16c6 0e0a      mvi      c, retry      ;max retries
224          16c8 c5          push     b
225          wboot0: ;enter here on error retries
226          16c9 010000      lxi      b, cpmb      ;set dma address to start of disk system
227          16cc cdbb17      call     setdma
228          16cf 0e00      mvi      c, 0          ;boot from drive 0
229          16d1 cd7d17      call     seldsk
230          16d4 0e00      mvi      c, 0
231          16d6 cda717      call     settrk      ;start with track 0
232          16d9 0e02      mvi      c, 2          ;start reading sector 2
233          16db cdac17      call     setsec

```

```

234      ;
235      ; read sectors, count nsects to zero
236      16de c1      pop    b      ;10-error count
237      16df 062c   mvi    b, nsects
238      rdsec:      ;read next sector
239      16e1 c5     push   b      ;save sector count
240      16e2 cdc117 call   read
241      16e5 c24917 jnz   booterr ;retry if errors occur
242      16e8 2a6c18 lhld  iod     ;increment dma address
243      16eb 118000 lxi   d, 128  ;sector size
244      16ee 19     dad   d      ;incremented dma address in hl
245      16ef 44     mov   b, h
246      16f0 4d     mov   c, l   ;ready for call to set dma
247      16f1 cdbb17 call   setdma
248      16f4 3a6b18 lda   ios     ;sector number just read
249      16f7 fe1a   cpi   26     ;read last sector?
250      16f9 da0517 jc    rd1
251      ; must be sector 26, zero and go to next track
252      16fc 3a6a18 lda   iot     ;get track to register a
253      16ff 3c     inr   a
254      1700 4f     mov   c, a   ;ready for call
255      1701 cda717 call   settrk
256      1704 af     xra   a      ;clear sector number
257      1705 3c     rd1: inr   a      ;to next sector
258      1706 4f     mov   c, a   ;ready for call
259      1707 cdac17 call   setsec
260      170a c1     pop   b      ;recall sector count
261      170b 05     dcr   b      ;done?
262      170c c2e116 jnz   rdsec
263      ;
264      ; done with the load, reset default buffer address
265      gocpm:      ;(enter here from cold start boot)
266      ; enable rst0 and rst7

```

```

!
267      170f f3          di
268      1710 3e12       mvi    a, 12h      ;initialize command
269      1712 d3fd       out    revrt
270      1714 af         xra    a
271      1715 d3fc       out    intc        ;cleared
272      1717 3e7e       mvi    a, inte     ;rst0 and rst7 bits on
273      1719 d3fc       out    intc
274      171b af         xra    a
275      171c d3f3       out    icon        ;interrupt control
276
277
278      171e 018000     lxi    b, buff
279      1721 cdbb17     call   setdma
280
281
282      1724 3ec3       mvi    a, jmp
283      1726 320000     sta    0
284      1729 210316     lxi    h, wboote
285      172c 220100     shld   1           ;jump wboot at location 00
286      172f 320500     sta    5
287      1732 210608     lxi    h, bdos
288      1735 220600     shld   6           ;jmp bdos at location 5
289
290      1738 323800     sta    7*8         ;jmp to mon80 (may have changed by ddt)
291      173b 2100f8     lxi    h, mon80
292      173e 223900     shld   7*8+1
293
294
295
296      1741 3a0400     lda    cdisk       ;last logged disk number
297      1744 4f         mov    c, a        ;send to ccp to log it in
298      1745 fb         ei
299      1746 c30000     jmp    cpmb
300

```

```
!
301 ; error condition occurred, print message and retry
302 booterr:
303 1749 c1 pop b ;recall counts
304 174a 0d dcr c
305 174b ca5217 jz booter0
306 ; try again
307 174e c5 push b
308 174f c3c916 jmp wboot0
309 ;
310 booter0:
311 ; otherwise too many retries
312 1752 215b17 lxi h, bootmsg
313 1755 cdd317 call prmsg
314 1758 c30fff jmp rmon80 ;mds hardware monitor
315 ;
316 bootmsg:
317 175b 3f626f6f74 db '?boot', 0
318 ;
319 ;
320 const: console status to reg-a
321 ; (exactly the same as mds call)
322 1761 c312f8 jmp csts
323 ;
324 conin: ;console character to reg-a
325 1764 cd03f8 call ci
326 1767 e67f ani 7fh ;remove parity bit
327 1769 c9 ret
328 ;
329 conout: ;console character from c to console out
330 176a c309f8 jmp co
331 ;
332 list: ;list device out
333 ; (exactly the same as mds call)
334 176d c30ff8 jmp lo
```

```

335 ;
336 listst:
337 ;return list status
338 1770 af xra a
339 1771 c9 ret ;always not ready
340 ;
341 punch: ;punch device out
342 ; (exactly the same as mds call)
343 1772 c30cf8 jmp po
344 ;
345 reader: ;reader character in to reg-a
346 ; (exactly the same as mds call)
347 1775 c306f8 jmp ri
348 ;
349 home: ;move to home position
350 ; treat as track 00 seek
351 1778 0e00 mvi c, 0
352 177a c3a717 jmp settrk
353 ;
354 seldsk: ;select disk given by register c
355 177d 210000 lxi h, 0000h ;return 0000 if error
356 1780 79 mov a, c
357 1781 fe04 cpi ndisks ;too large?
358 1783 d0 rnc ;leave hl = 0000
359 ;
360 1784 e602 ani 10b ;00 00 for drive 0, 1 and 10 10 for drive 2, 3
361 1786 326618 sta dbank ;to select drive bank
362 1789 79 mov a, c ;00, 01, 10, 11
363 178a e601 ani 1b ;mds has 0, 1 at 78, 2, 3 at 88
364 178c b7 ora a ;result 00?
365 178d ca9217 jz setdrive
366 1790 3e30 mvi a, 00110000b ;selects drive 1 in bank
367 setdrive:
368 1792 47 mov b, a ;save the function

```

```

369      1793 216818      lxi   h, iof      ;io function
370      1796 7e          mov   a, m
371      1797 e6cf        ani   11001111b   ;mask out disk number
372      1799 b0          ora   b            ;mask in new disk number
373      179a 77          mov   m, a        ;save it in iopb
374      179b 69          mov   l, c
375      179c 2600        mvi   h, 0        ;hl=disk number
376      179e 29          dad   h            ;*2
377      179f 29          dad   h            ;*4
378      17a0 29          dad   h            ;*8
379      17a1 29          dad   h            ;*16
380      17a2 113316      lxi   d, dpbase
381      17a5 19          dad   d            ;hl=disk header table address
382      17a6 c9          ret
383      ;
384      ;
385      ;settrk:         ;set track address given by c
386      17a7 216a18      lxi   h, iot
387      17aa 71          mov   m, c
388      17ab c9          ret
389      ;
390      ;setsec:         ;set sector number given by c
391      17ac 216b18      lxi   h, ios
392      17af 71          mov   m, c
393      17b0 c9          ret
394      ;sectran:
395      ;translate sector bc using table at de
396      17b1 0600        mvi   b, 0        ;double precision sector number in bc
397      17b3 eb          xchg                    ;translate table address to hl
398      17b4 09          dad   b            ;translate (sector) address
399      17b5 7e          mov   a, m        ;translated sector number to a
400      17b6 326b18      sta   ios
401      17b9 6f          mov   l, a        ;return sector number in l
402      17ba c9          ret

```

```

403      ;
404      setdma: ;set dma address given by regs b, c
405      17bb 69      mov    l, c
406      17bc 60      mov    h, b
407      17bd 226c18  shld  iod
408      17c0 c9      ret
409      ;
410      read: ;read next disk record (assuming disk/trk/ sec/dma set)
411      17c1 0e04      mvi   c, readf ;set to read function
412      17c3 cde017    call  setfunc
413      17c6 cdf017    call  waitio ;perform read function
414      17c9 c9      ret ;may have error set in reg-a
415      ;
416      ;
417      write: ;disk write function
418      17ca 0e06      mvi   c, writf
419      17cc cde017    call  setfunc ;set to write function
420      17cf cdf017    call  waitio
421      17d2 c9      ret ;may have error set
422      ;
423      ;
424      ; utility subroutines
425      prmsg: ;print message at h, l to 0
426      17d3 7e      mov   a, m
427      17d4 b7      ora   a      zero?
428      17d5 c8      rz
429      ; more to print
430      17d6 e5      push  h
431      17d7 4f      mov   c, a
432      17d8 cd6a17   call  conout
433      17db e1      pop   h
434      17dc 23      inc   h
435      17dd c3d317   jmp   prmsg
436      ;

```

```

437                                     setfunc:
438                                     ;
439      17e0 216818                       lxi   h, iof       ;io function address
440      17e3 7e                           mov   a, m         ;get it to accumulator for masking
441      17e4 e6f8                           ani   11111000b    ;remove previous command
442      17e6 b1                             ora   c            ;set to new command
443      17e7 77                             mov   m, a        ;replaced in iopb
444                                     ;
445                                     ; the mds-800 controller requires disk bank bit in sector byte
446                                     ; mask the bit from the current i/o function
447      17e8 e620                           ani   00100000b    ;mask the disk select bit
448      17ea 216b18                        lxi   h, ios       ;address the sector select byte
449      17ed b6                             ora   m            ;select proper disk bank
450      17ee 77                             mov   m, a        ;set disk select bit on/off
451      17ef c9                             ret
452                                     ;
453      17f0 0e0a                           mvi   c, retry     ;max retries before perm error
454                                     ;
455      17f2 cd3f18                        call  intype        ;in rtype
456      17f5 cd4c18                        call  inbyte        ;clears the controller
457                                     ;
458                                     ;
459      17f8 3a6618                        lda   dbank         ;set bank flags
460      17fb b7                             ora   a            ;zero if drive 0, 1 and nz if 2, 3
461      17fc 3e67                           mvi   a, iopb and offh ;low address for iopb
462      17fe 0618                           mvi   b, iopb shr 8 ;high address for iopb
463      1800 c20b18                        jnz   iod1         ;drive bank 1?
464      1803 d379                           out   ilow          ;low address to controller
465      1805 78                             mov   a, b
466      1806 d37a                           out   ihigh         ;high address
467      1808 c31018                        jmp   waito        ;to wait for complete
468                                     ;
469      180b d389                           iod1: ;drive bank 1
470                                     out   ilow+10h     ;88 for drive bank 10

```

```

471      180d 78          mov   a, b
472      180e d38a       out   ihigh+10h
473      ;
474      1810 cd5918     waito: call  instat          ;wait for completion
475      1813 e604       ani   iordy          ;ready?
476      1815 ca1018     jz    waito
477      ;
478      ;
479      1818 cd3f18     call  intype          ;must be io complete (00) unlinked
480      ;               00 unlinked i/o complete, 01 linked i/o complete (not used)
481      ;               io disk status changed 11 (not used)
482      181b fe02       cpi   10b            ;ready status change?
483      181d ca3218     jz    wready
484      ;
485      ;               must be 00 in the accumulator
486      1820 b7         ora   a
487      1821 c23818     jnz   werror          ;some other condition, retry
488      ;
489      ;               check i/o error bits
490      1824 cd4c18     call  inbyte
491      1827 17         ral
492      1828 da3218     jc    wready          ;unit not ready
493      182b 1f         rar
494      182c e6fe       ani   11111110b      ;any other errors? (deleted data ok)
495      182e c23818     jnz   werror
496      ;
497      ;               read or write is ok, accumulator contains zero
498      1831 c9         ret
499      ;
500      wready: ;not ready, treat as error for now
501      1832 cd4c18     call  inbyte          ;clear result byte
502      1835 c33818     jmp   trycount
503      ;
504      werror: ;return hardware malfunction (crc, track, seek, etc.)

```

```

505 ; the mds controller has returned a bit in each position
506 ; of the accumulator, corresponding to the conditions:
507 ; 0 —deleted data (accepted as ok above)
508 ; 1 —crc error
509 ; 2 —seek error
510 ; 3 —address error (hardware malfunction)
511 ; 4 —data over/under flow (hardware malfunction)
512 ; 5 —write protect (treated as not ready)
513 ; 6 —write error (hardware malfunction)
514 ; j —not ready
515 ; (accumulator bits are numbered 7 6 5 4 3 2 1 0)
516 ;
517 ; it may be useful to filter out the various conditions,
518 ; but we will get a permanent error message if it is not
519 ; recoverable. in any case, the not ready condition is
520 ; treated as a separated condition for later improvement
521 trycount:
522 ; register c contains retry count, decrement 'til zero
523 1838 0d dcr c
524 1839 c2f217 jnz rewait ;for another try
525 ;
526 ; cannot recover from error
527 183c 3e01 mvi a, 1 ;error code
528 183e c9 ret
529 ;
530 ; intype, inbyte, instat read drive bank 00 or 10
531 183f 3a6618 intype: lda dbank
532 1842 b7 ora a
533 1843 c24918 jnz intyp1 ;skip to bank 10
534 1846 db79 in rtype
535 1848 c9 ret
536 1849 db89 intyp1: in rtype+10h ;78 for 0, 1 88 for 2, 3
537 184b c9 ret

```

```

538      ;
539      184c 3a6618      inbyte:  lda    dbank
540      184f b7          ora    a
541      1850 c25618      jnz    inbyt1
542      1853 db7b        in    rbyte
543      1855 c9          ret
544      1856 db8b        inbyt1: in    rbyte+10h
545      1858 c9          ret
546      ;
547      1859 3a6618      instat: lda    dbank
548      185c b7          ora    a
549      185d c26318      jnz    insta1
550      1860 db78        in    dstat
551      1862 c9          ret
552      1863 db88        insta1: in    dstat+10h
553      1865 c9          ret
554      ;
555      ;
556      ;
557      ;      data areas (must be in ram)
558      1866 00          dbank:  db    0          ;disk bank 00 if drive 0, 1
559      ;              ;              10 if drive 2, 3
560      iopb:           ;io parameter block
561      1867 80          db    80h          ;normal i/o operation
562      1868 04          iof:   db    readf    ;io function, initial read
563      1869 01          ion:   db    1          ;number of sectors to read
564      186a 02          iot:   db    offset   ;track number
565      186b 01          ios:   db    1          ;sector number
566      186c 8000        iod:   dw    buff     ;io address
567      ;
568      ;
569      ;      define ram areas for bdos operation
570      ;      endif

```

```

571      186e+=      begdat  equ  $
572      186e+      dirbuf: ds  128      ;directory access buffer
573      18ee+      alv0:   ds  31
574      190d+      csv0:   ds  16
575      191d+      alv1:   ds  31
576      193c+      csv1:   ds  16
577      194c+      alv2:   ds  31
578      196b+      csv2:   ds  16
579      197b+      alv3:   ds  31
580      199a+      csv3:   ds  16
581      19aa+=     enddat  equ  $
582      013c+=     datsiz  equ  $-begdat
583      19aa      end

```

```

als1      001f      141#
als2      001f      146#
als3      001f      151#
alv0      18ee      87      573#
alv1      191d      91      575#
alv2      194c      95      577#
alv3      197b      99      579#
base      0078      180#      181      182      183      185      186
bdos      0806      29#      287
begdat    186e      571#      582
bias      0000      19#      22#
boot      16b3      63      207#
booter0   1752      305      310#
booterr   1749      241      302#
`bootmsg  175b      312      316#
buff      0080      34#      209      221      278      566
cdisk     0004      33#      213      296
ci        f803      172#      325
co        f809      174#      330

```

conin	1764	66	324#					
conout	176a	67	329#	432				
const	1761	65	320#					
cpmb	0000	28#	29	30	226	299		
cpml	1600	30#	31					
cr	000d	192#	196	205				
css1	0010	142#						
css2	0010	147#						
css3	0010	152#						
csts	f812	177#	322					
csv0	190d	87	574#					
csv1	193c	91	576#					
csv2	196b	95	578#					
csv3	199a	99	580#					
datsiz	013c	582#						
dbank	1866	361	459	531	539	547	558#	
dirbuf	186e	86	90	94	98	572#		
dpb0	1673	86	101#	140	145	150		
dpb1	1673	90	140#					
dpb2	1673	94	145#					
dpb3	1673	98	150#					
dpbase	1633	83#	380					
dpe0	1633	84#						
dpe1	1643	88#						
dpe2	1653	92#						
dpe3	1663	96#						
dstat	0078	181#	550	552				
enddat	19aa	581#						
false	0000	15#	16					
gocpm	170f	214	265#					
home	1778	71	349#					
icon	00f3	166#	275					
ihigh	007a	186#	466	472				

ilow	0079	185#	464	470		
inbyt1	1856	541	544#			
inbyte	184c	457	490	501	539#	
insta1	1863	549	552#			
instat	1859	474	547#			
intc	00fc	165#	271	273		
inte	007e	167#	272			
intyp1	1849	533	536#			
intype	183f	456	479	531#		
iod	186c	242	407	566#		
iodr1	180b	463	469#			
iof	1868	369	439	562#		
ion	1869	563#				
iopb	1867	461	462	560#		
iordy	0004	191#	475			
ios	186b	248	391	400	447	565#
iot	186a	252	386	564#		
lf	000a	193#	196	196	205	
list	176d	68	332#			
listst	1770	78	336#			
lo	f80f	176#	334			
mon80	f800	170#	291			
nsects	002c	31#	237			
offset	0002	32#	100	564		
patch	1600	25#	27	28		
po	f80c	175#	343			
prmsg	17d3	211	313	425#	435	
punch	1772	69	341#			
rbyte	007b	183#	542	544		
rd1	1705	250	257#			
rdsec	16e1	238#	262			
read	17c1	76	240	410#		
reader	1775	70	345#			
readf	0004	188#	411	562		

recal	0003	190#					
retry	000a	35#	223	453			
revrt	00fd	164#	269				
rewait	17f2	454#	524				
ri	f806	173#	347				
rmon80	ff0f	171#	314				
rtype	0079	182#	534	536			
sectran	17b1	79	394#				
seldsk	177d	72	229	354#			
setdma	17bb	75	227	247	279	404#	
setdrive	1792	365	367#				
setfunc	17e0	412	419	437#			
setsec	17ac	74	233	259	390#		
settrk	17a7	73	231	255	352	385#	
signon	169c	195#	210				
test	0000	16#	18	21	197	200	289
true	ffff	14#	15				
trycount	1838	502	521#				
vers	0016	6#	204	204			
waito	1810	467	474#	476			
waitio	17f0	413	420	452#			
wboot	16c3	64	217#				
wboot0	16c9	225#	308				
wboote	1603	64#	284				
werror	1838	487	495	504#			
wready	1832	483	492	500#			
write	17ca	77	417#				
writf	0006	189#	418				
xlt0	1682	84	112#	143	148	153	
xlt1	1682	88	143#				
xlt2	1682	92	148#				
xlt3	1682	96	153#				



Appendix B: A Skeletal CBIOS

```

1          ;          skeletal cbios for first level of cp/m 2.0 alteration
2          ;
3      0014 =      msize  equ   20          ;cp/m version memory size in kilobytes
4          ;
5          ;          "bias" is address offset from 3400h for memory systems
6          ;          than 16k (referred to as "b" throughout the text)
7          ;
8      0000 =      bias   equ   (msize-20)*1024
9      3400 =      ccp    equ   3400h+bias  ;base of ccp
10     3c06 =      bdos   equ   ccp+806h   ;base of bdos
11     4a00 =      bios   equ   ccp+1600h  ;base of bios
12     0004 =      cdisk  equ   0004h     ;current disk number 0=a, . . . , 15=p
13     0003 =      iobyte equ   0003h     ;intel i/o byte
14          ;
15     4a00          org   bios           ;origin of this program
16     002c =      nsects equ   ($-ccp)/128 ;warm start sector count
17          ;
18          ;          jump vector for individual subroutines
19     4a00 c39c4a   jmp   boot           ;cold start
20     4a03 c3a64a   wboote: jmp   wboot        ;warm start
21     4a06 c3114b   jmp   const          ;console status
22     4a09 c3244b   jmp   conin         ;console character in
23     4a0c c3374b   jmp   conout        ;console character out
24     4a0f c3494b   jmp   list          ;list character out
25     4a12 c34d4b   jmp   punch        ;punch character out
26     4a15 c34f4b   jmp   reader        ;reader character out
27     4a18 c3544b   jmp   home         ;move head to home position

```

```

28 4a1b c35a4b jmp seldsk ;select disk
29 4a1e c37d4b jmp settrk ;set track number
30 4a21 c3924b jmp setsec ;set sector number
31 4a24 c3ad4b jmp setdma ;set dma address
32 4a27 c3c34b jmp read ;read disk
33 4a2a c3d64b jmp write ;write disk
34 4a2d c34b4b jmp listst ;return list status
35 4a30 c3a74b jmp sectran ;sector translate
36 ;
37 ; fixed data tables for four-drive standard
38 ; ibm-compatible 8" disks
39 ; disk parameter header for disk 00
40 4a33 734a0000 dpbase: dw trans, 0000h
41 4a37 00000000 dw 0000h, 0000h
42 4a3b f04c8d4a dw dirbf, dpblk
43 4a3f ec4d704d dw chk00, all00
44 ; disk parameter header for disk 01
45 4a43 734a0000 dw trans, 0000h
46 4a47 00000000 dw 0000h, 0000h
47 4a4b f04c8d4a dw dirbf, dpblk
48 4a4f fc4d8f4d dw chk01, all01
49 ; disk parameter header for disk 02
50 4a53 734a0000 dw trans, 0000h
51 4a57 00000000 dw 0000h, 0000h
52 4a5b f04c8d4a dw dirbf, dpblk
53 4a5f 0c4eae4d dw chk02, all02
54 ; disk parameter header for disk 03
55 4a63 734a0000 dw trans, 0000h
56 4a67 00000000 dw 0000h, 0000h
57 4a6b f04c8d4a dw dirbf, dpblk
58 4a6f 1c4ecd4d dw chk03, all03
59 ;
60 ; sector translate vector

```

```

61      4a73 01070d13      trans:  db    1, 7, 13, 19    ;sectors 1, 2, 3, 4
62      4a77 19050b11      db    25, 5, 11, 17    ;sectors 5, 6, 7, 8
63      4a7b 1703090f      db    23, 3, 9, 15    ;sectors 9, 10, 11, 12
64      4a7f 1502080e      db    21, 2, 8, 14    ;sectors 13, 14, 15, 16
65      4a83 141a060c      db    20, 26, 6, 12   ;sectors 17, 18, 19, 20
66      4a87 1218040a      db    18, 24, 4, 10   ;sectors 21, 22, 23, 24
67      4a8b 1016          db    16, 22          ;sectors 25, 26
68
69      ;
70      4a8d 1a00          dpblk: ;disk parameter block, common to all disks
71      4a8f 03          dw    26              ;sectors per track
72      4a90 07          db    3                ;block shift factor
73      4a91 00          db    7                ;block mask
74      4a92 f200        dw    0                ;null mask
75      4a94 3f00        dw    242             ;disk size-1
76      4a96 c0          dw    63              ;directory max
77      4a97 00          db    192             ;alloc 0
78      4a98 1000        db    0                ;alloc 1
79      4a9a 0200        dw    16              ;check size
80      ;                dw    2                ;track offset
81      ;
82      ;                end of fixed tables
83      ;
84      ;                individual subroutines to perform each function
85      4a9c af          boot: ;simplest case is to just perform parameter initialization
86      4a9d 320300      xra   a                ;zero in the accum
87      4aa0 320400      sta  iobyte            ;clear the iobyte
88      4aa3 c3ef4a      sta  cdisk             ;select disk zero
89      ;                jmp  gocpm              ;initialize and go to cp/m
90      ;
91      4aa6 318000      wboot: ;simplest case is to read the disk until all sectors loaded
92      4aa9 0e00        lxi  sp, 80h           ;use space below buffer for stack
93      4aab cd5a4b      mvi  c, 0              ;select disk 0
94      4aae cd544b      call seldsk           ;go to track 00

```

```

95      ;
96      4ab1 062c      mvi    b, nsects    ;b counts # of sectors to load
97      4ab3 0e00      mvi    c, 0         ;c has the current track number
98      4ab5 1602      mvi    d, 2         ;d has the next sector to read
99      ;
100     ;              note that we begin by reading track 0, sector 2 since sector 1
101     4ab7 210034    lxi    h, ccp       ;base of cp/m (initial load point)
102     load1:         ;load one more sector
103     4aba c5        push   b           ;save sector count, current track
104     4abb d5        push   d           ;save next sector to read
105     4abc e5        push   h           ;save dma address
106     4abd 4a        mov    c, d        ;get sector address to register c
107     4abe cd924b    call  setsec      ;set sector address from register c
108     4ac1 c1        pop    b           ;recall dma address to b, c
109     4ac2 c5        push   b           ;replace on stack for later recall
110     4ac3 cdad4b    call  setdma      ;set dma address from b, c
111     ;
112     ;              drive set to 0, track set, sector set, dma address set
113     4ac6 cdc34b    call  read
114     4ac9 fe00      cpi    00h         ;any errors?
115     4acb c2a64a    jnz   wboot       ;retry the entire boot if an error occurs
116     ;
117     ;              no error, move to next sector
118     4ace e1        pop    h           ;recall dma address
119     4acf 118000    lxi    d, 128     ;dma=dma+128
120     4ad2 19        dad    d           ;new dma address is in h, l
121     4ad3 d1        pop    d           ;recall sector address
122     4ad4 c1        pop    b           ;recall number of sectors remaining, and current trk
123     4ad5 05        dcr   b           ;sectors=sectors-1
124     4ad6 caef4a    jz    gocpm       ;transfer to cp/m if all have been loaded
125     ;
126     ;              more sectors remain to load, check for track change
127     4ad9 14        inr   d

```

```

128      4ada 7a      mov  a, d      ;sector=27?, if so, change tracks
129      4adb fe1b   cpi  27
130      4add daba4a jc   load1     ;carry generated if sector<27
131      ;
132      ;          end of current track, go to next track
133      4ae0 1601   mvi  d, 1     ;begin with first sector of next track
134      4ae2 0c     inr  c        ;track=track+1
135      ;
136      ;          save register state, and change tracks
137      4ae3 c5     push b
138      4ae4 d5     push d
139      4ae5 e5     push h
140      4ae6 cd7d4b call settrk   ;track address set from register c
141      4ae9 e1     pop  h
142      4aea d1     pop  d
143      4aeb c1     pop  b
144      4aec c3ba4a jmp  load1     ;for another sector
145      ;
146      ;          end of load operation, set parameters and go to cp/m
147      ;          gocpm:
148      4aef 3ec3   mvi  a, 0c3h  ;c3 is a jmp instruction
149      4af1 320000 sta  0        ;for jmp to wboot
150      4af4 21034a lxi  h, wboote ;wboot entry point
151      4af7 220100 shld 1        ;set address field for jmp at 0
152      ;
153      4afa 320500 sta  5        ;for jmp to bdos
154      4afd 21063c lxi  h, bdos  ;bdos entry point
155      4b00 220600 shld 6        ;address field of jump at 5 to bdos
156      ;
157      4b03 018000 lxi  b, 80h   ;default dma address is 80h
158      4b06 cdad4b call setdma
159      ;
160      4b09 fb     ei          ;enable the interrupt system

```

```

161      4b0a 3a0400      lda    cdisk      ;get current disk number
162      4b0d 4f          mov    c, a       ;send to the ccp
163      4b0e c30034     jmp    ccp        ;go to cp/m for further processing
164
165      ;
166      ;              simple i/o handlers (must be filled in by user)
167      ;              in each case, the entry point is provided, with space reserved
168      ;              to insert your own code
169      ;
170      const:         ;console status, return 0ffh if character ready, 00h if not
171      4b11           ds     10h        ;space for status subroutine
172      4b21 3e00     mvi    a, 00h
173      4b23 c9       ret
174      ;
175      conin:         ;console character into register a
176      4b24           ds     10h        ;space for input routine
177      4b34 e67f     ani    7fh        ;strip parity bit
178      4b36 c9       ret
179      ;
180      conout:        ;console character output from register c
181      4b37 79       mov    a, c       ;get to accumulator
182      4b38           ds     10h        ;space for output routine
183      4b48 c9       ret
184      ;
185      list:          ;list character from register c
186      4b49 79       mov    a, c       ;character to register a
187      4b4a c9       ret              ;null subroutine
188      ;
189      listst:        ;return list status (0 if not ready, 1 if ready)
190      4b4b af       xra    a          ;0 is always ok to return
191      4b4c c9       ret
192      ;
193      punch:         ;punch character from register c

```

```

194      4b4d 79          mov   a, c          ;character to register a
195      4b4e c9          ret                   ;null subroutine
196      ;
197      ;
198      reader:        ;read character into register a from reader device
199      4b4f 3e1a        mvi   a, 1ah        ;enter end of file for now (replace later)
200      4b51 e67f        ani   7fh           ;remember to strip parity bit
201      4b53 c9          ret
202      ;
203      ;
204      ;              i/o drivers for the disk follow
205      ;              for now, we will simply store the parameters away for use
206      ;              in the read and write subroutines
207      ;
208      home:          ;move to the track 00 position of current drive
209      ;              translate this call into a settrk call with parameter 00
210      4b54 0e00        mvi   c, 0          ;select track 0
211      4b56 cd7d4b      call  settrk
212      4b59 c9          ret                   ;we will move to 00 on first read/write
213      ;
214      seldsk:        ;select disk given by register c
215      4b5a 210000      lxi   h, 0000h      ;error return code
216      4b5d 79          mov   a,c.
217      4b5e 32ef4c      sta   diskno
218      4b61 fe04        cpi   4              ;must be between 0 and 3
219      4b63 d0          rnc                   ;no carry if 4, 5, . . .
220      ;              disk number is in the proper range
221      4b64             ds    10              ;space for disk select
222      ;              compute proper disk parameter header address
223      4b6e 3aef4c      lda   diskno
224      4b71 6f          mov   l, a           ;l=disk number 0, 1, 2, 3
225      4b72 2600        mvi   h, 0          ;high order zero
226      4b74 29          dad   h              ;*2

```

```

227 4b75 29      dad  h      ;*4
228 4b76 29      dad  h      ;*8
229 4b77 29      dad  h      ;*16 (size of each header)
230 4b78 11334a  lxi  d, dpbase
231 4b7b 19      dad  0      ;hl=.dpbase(diskno*16)
232 4b7c c9      ret
233
234          ;
settrk:      ;set track given by register c
235 4b7d 79      mov  a, c
236 4b7e 32e94c  sta  track
237 4b81          ds   10h      ;space for track select
238 4b91 c9      ret
239
240          ;
setsec:      ;set sector given by register c
241 4b92 79      mov  a, c
242 4b93 32eb4c  sta  sector
243 4b96          ds   10h      ;space for sector select
244 4ba6 c9      ret
245
246          ;
sectran:     ;translate the sector given by bc using the
247          ;translate table given by de
248          xchg          ;hl=.trans
249 4ba7 eb      dad  b      ;hl=.trans(sector)
250 4ba8 09      mov  l, m   ;l = trans(sector)
251 4ba9 6e      mvi  h, 0   ;hl = trans(sector)
252 4baa 2600    ret        ;with value in hl
253 4bac c9
254          ;
setdma:      ;set dma address given by registers b and c
255          mov  l, c      ;low order address
256 4bad 69      mov  h, b   ;high order address
257 4bae 60      shld dmaad ;save the address
258 4baf 22ed4c ds   10h    ;space for setting the dma address
259 4bb2

```

```

260      4bc2 c9          ret
261          ;
262          read:      ;perform read operation (usually this is similar to write
263          ;          so we will allow space to set up read command, then use
264          ;          common code in write)
265      4bc3          ds    10h          ;set up read command
266      4bd3 c3e64b    jmp    waitio      ;to perform the actual i/o
267          ;
268          write:     ;perform a write operation
269      4bd6          ds    10h          ;set up write command
270          ;
271          waitio:    ;enter here from read and write to perform the actual i/o
272          ;          operation. return a 00h in register a if the operation completes
273          ;          properly, and 01h if an error occurs during the read or write
274          ;
275          ;          in this case, we have saved the disk number in 'diskno' (0, 1)
276          ;          the track number in 'track' (0-76)
277          ;          the sector number in 'sector' (1-26)
278          ;          the dma address in 'dmaad' (0-65535)
279      4be6          ds    256         ;space reserved for i/o drivers
280      4ce6 3e01      mvi    a, 1        ;error condition
281      4ce8 c9          ret          ;replaced when filled- in
282          ;
283          ;          the remainder of the cbios is reserved uninitialized
284          ;          data area, and does not need to be a part of the
285          ;          system memory image (the space must be available,
286          ;          however, between "begdat" and "enddat").
287          ;
288      4ce9          track:   ds    2          ;two bytes for expansion
289      4ceb          sector:  ds    2          ;two bytes for expansion
290      4ced          dmaad:   ds    2          ;direct memory address

```

```

291     4cef          diskno:  ds    1          ;disk number 0-15
292          ;
293          ;          scratch ram area for bdos use
294     4cf0 =       begdat   equ    $          ;beginning of data area
295     4cf0         dirbf:   ds    128        ;scratch directory area
296     4d70         all00:   ds    31         ;allocation vector 0
297     4d8f         all01:   ds    31         ;allocation vector 1
298     4dae         all02:   ds    31         ;allocation vector 2
299     4dcd         all03:   ds    31         ;allocation vector 3
300     4dec         chk00:   ds    16         ;check vector 0
301     4dfc         chk01:   ds    16         ;check vector 1
302     4e0c         chk02:   ds    16         ;check vector 2
303     4e1c         chk03:   ds    16         ;check vector 3
304          ;
305     4e2c         enddat   equ    $          ;end of data area
306     013c =       datsiz   equ    $-begdat; ;size of data area
307     4e2c         end

```

```

all00      4d70      43      296#
all01      4d8f      48      297#
all02      4dae      53      298#
all03      4dcd      58      299#
bdos       3c06      10#     154
begdat     4cf0      294#     306
bias       0000      8#       9
bios       4a00      11#      15
boot       4a9c      19       84#
ccp        3400      9#       10      11      16      101      163
cdisk      0004      12#      87      161
chk00      4dec      43       300#
chk01      4dfc      48       301#
chk02      4e0c      53       302#
chk03      4e1c      58       303#

```

conin	4b24	22	175#			
conout	4b37	23	180#			
const	4b11	21	170#			
datsiz	013c	306#				
dirbf	4cf0	42	47	52	57	295#
diskno	4cef	217	223	291#		
dmaad	4ced	258	290#			
dpbase	4a33	40#	230			
dpblk	4a8d	42	47	52	57	69#
enddat	4e2c	305#				
gocpm	4aef	88	124	147#		
home	4b54	27	94	208#		
iobyte	0003	13#	86			
list	4b49	24	185#			
listst	4b4b	34	189#			
load1	4aba	102#	130	144		
msize	0014	3#	8			
nsects	002c	16#	96			
punch	4b4d	25	193#			
read	4bc3	32	113	262#		
reader	4b4f	26	198#			
sector	4ceb	242	289#			
sectran	4ba7	35	246#			
seldsk	4b5a	28	93	214#		
setdma	4bad	31	110	158	255#	
setsec	4b92	30	107	240#		
settrk	4b7d	29	140	211	234#	
track	4ce9	236	288#			
trans	4a73	40	45	50	55	61#
waitio	4be6	266	271#			
wboot	4aa6	20	90#	115		
wboote	4a03	20#	150			
write	4bd6	33	268#			



Appendix C: A Skeletal GETSYS/PUTSYS Program

```

:
:
: combined getsys and putsys programs from
: Sec 6.4
:
: Start the programs at the base of the TPA
0100      org 0100h

0014 =    msize      equ 20          ; size of cp/m in Kbytes

: "bias" is the amount to add to addresses for > 20k
: (referred to as "b" throughout the text)
0000 =    bias       equ (msize-20)*1024
3400 =    ccp        equ 3400h+bias
3c00 =    bdos       equ ccp+0800h
4a00 =    bios       equ ccp+1600h

:
: getsys programs tracks 0 and 1 to memory at
: 3880h + bias
:
: register      usage
: a             (scratch register)
: b             track count (0...76)
: c             sector count (1...26)
: d,e          (scratch register pair)
: h,l          load address
: sp           set to track address

gstart:
lxi      sp,ccp-0080h      ; start of getsys
lxi      h,ccp-0080h      ; convenient place
mvi     b,0               ; set initial load
: start with track
rd$trk: ; read next track
0100 318033
0103 218033
0106 0600

```

```

0108 0e01          mvi c,1          ; each track start

rd$sec:
010a cd0003      call rd$sec          ; get the next sector
010d 118000      lxi d,128          ; offset by one sector
0110 19          dad d              ; (hl=hl+128)
0111 0c          inr c              ; next sector
0112 79          mov a,c            ; fetch sector number
0113 felb       cpi 27          ; and see if last
0115 da0a01      jc rdsec          ; <, do one more

; arrive here at end of track, move to next track

0118 04          inr b              ; track = track+1
0119 78          mov a,b            ; check for last
011a fe02       cpi 2              ; track = 2 ?
011c da0801      jc rd$trk         ; <, do another

; arrive here at end of load, halt for lack of anything
; better
ei
hlt
;
; putsys program, places memory image
; starting at
; 3880h + bias back to tracks 0 and 1
; start this program at the next page boundary
; org ($+0100h) and 0ff00h

0200          putsys:
0200 318033      lxi sp,ccp-0080h  ; convenient place
0203 218033      lxi h,ccp-0080h  ; start of dump
0206 0600       mvi b,0          ; start with track

wr$trk:
0208 0e01       mvi c,1          ; start with sector

wr$sec:
020a cd0004      call wr$sec        ; write one sector
020d 118000      lxi d,128          ; length of each
0210 19          dad d              ; <hl>=<hl> + 128
0211 0c          inr c              ; <c>=<c> + 1
0212 79          mov a,c            ; see if
0213 felb       cpi 27          ; past end of track
0215 da0a02      jc wr$sec        ; no, do another

; arrive here at end of track, move to next track

0218 04          inr b              ; track = track+1
0219 78          mov a,b            ; see if
021a fe02       cpi 2              ; last track
021c da0802      jc wr$trk         ; no, do another

;
; done with putsys, halt for lack of anything
; better
ei
hit
021f fb
0220 76

```

```

; user supplied subroutines for sector read and write
;
;      move to next page boundary
0300      org ($+0100h) and 0ff00h

read$sec:
; read the next sector
; track in <b>,
; sector in <c>
; dmaaddr in <h>
0300 c5      pushb
0301 e5      pushh

; user defined read operation goes here
0302      ds 64

0342 e1      pop h
0343 c1      pop b
0344 c9      ret

0400      org ($+0100h) and 0ff00h ;another page
; boundary
write$sec:
; same parameters as read$sec
0400 c5      pushb
0401 e5      pushh

; user defined write operation goes here
0402      ds 64

0442 e1      pop h
0443 c1      pop b
0444 c9      ret

; end of getsys/putsys program
0445      end

```



Appendix D: The MDS-800 Cold Start Loader for CP/M 2

```

1          title   mds cold start loader at 3000h'
2          ;
3          mds-800 cold start loader for cp/m 2.0
4          ;
5          ;
6          ;
7          0000 =   false   equ   0
8          ffff   true    equ   not false
9          0000 =   testing equ   false    if true, then go to mon80 on errors
10         ;
11         ;
12         bias    equ   03400h
13         endif
14         ;
15         0000 =   bias    equ   0000h
16         endif
17         cpmb    equ   bias      ;base of dos load
18         bdos    equ   806h+bias ;entry to dos for calls
19         bdose   equ   1880h+bias ;end of dos load
20         boot    equ   1600h+bias ;cold start entry point
21         rboot   equ   boot+3    ;warm start entry point
22         ;
23         3000    org   03000h    ;loaded down from hardware boot at 3000H
24         ;
25         1880 =   bdosl   equ   bdose-cpmb
26         0002 =   ntrks   equ   2      ;number of tracks to read
27         0031 =   bdoss   equ   bdosl/128 ;number of sectors in dos
28         0019 =   bdoso   equ   25     ;number of bdos sectors on track 0
29         0018 =   bdos1   equ   bdoss-bdoso ;number of sectors on track 1

```

```

30      ;
31      f800 =      mon80 equ    of800h ;intel monitor base
32      ffof =      rmon80 equ    offofh ;restart location for mon80
33      0078 =      base  equ    078h  ;'base' used by controller
34      0079 =      rtype equ    base+1 ;result type
35      007b =      rbyte equ    base+3 ;result byte
36      007f =      reset equ    base+7 ;reset controller
37      ;
38      0078 =      dstat equ    base   ;disk status port
39      0079 =      ilow  equ    base+1 ;low iopb address
40      007a =      ihigh equ    base+2 ;high iopb address
41      00ff =      bsw   equ    offh   ;boot switch
42      0003 =      recal equ    3h     ;recalibrate selected drive
43      0004 =      readf equ    4h     ;disk read function
44      0100 =      stack equ    100h   ;use end of boot for stack
45      ;
46      rstart:
47      3000 310001      lxi    sp,stack; ;in case of call to mon80
48      ;                clear disk status
49      3003 db79        in     rtype
50      3005 db7b        in     rbyte
51      ;                check if boot switch is off
52      coldstart:
53      3007 dbff        in     bsw
54      3009 e602        ani    02h     ;switch on?
55      300b c20730      jnz    coldstart
56      ;                clear the controller
57      300e d37f        out    reset   ;logic cleared
58      ;
59      ;
60      3010 0602        mvi    b,ntrks ;number of tracks to read
61      3012 214230      lxi    h,iopbo
62      ;

```

```

63          start:
64          ;
65          ;      read first/next track into cpmb
66          3015 7d      mov    a,l
67          3016 d379    out    ilow
68          3018 7c      mov    a,h
69          3019 d37a    out    ihigh
70          301b db78    waito:  in    dstat
71          301d e604    ani    4
72          301f ca1b30  jz     waito
73          ;
74          ;      check disk status
75          3022 db79    in     rtype
76          3024 e603    ani    11b
77          3026 fe02    cpi    2
78          ;
79          ;      if testing
80          cnc    rmon80 ;go to monitor if 11 or 10
81          endif
82          ;      if not testing
83          3028 d20030  jnc    rstart ;retry the load
84          endif
85          ;
86          302b db7b    in     rbyte ;i/o complete, check status
87          ;      if not ready, then go to mon80
88          302d 17      ral
89          302e dc0fff  cc     rmon80 ;not ready bit set
90          3031 1f      rar    ;restore
91          3032 e61e    ani    11110b ;overrun/addr err/seek/crc/xxxx
92          ;
93          ;      if testing
94          cnz    rmon80 ;go to monitor
95          endif
96          ;      if not testing

```

```

97      3034 c20030      jnz  rstart      ;retry the load
98      ;
99      ;
100     ;
101     3037 110700      lxi  d,iopbl     ;length of iopb
102     303a 19          dad  d           ;addressing next iopb
103     303b 05          dcr  b           ;count down tracks
104     303c c21530      jnz  start
105     ;
106     ;
107     ;
108     303f c30016      jmp  to boot to print initial message, and set up jmps
109     ;
110     ;
111     3042 80          iopbo: db 80h     ;iocw, no update
112     3043 04          db  readf       ;read function
113     3044 19          db  bdoso       ;# sectors to read on track 0
114     3045 00          db  0           ;track 0
115     3046 02          db  2           ;start with sector 2 on track 0
116     3047 0000       dw  cpmb        ;start at base of bdos
117     0007 =          iopbl equ $-iopbo
118     ;
119     3049 80          iopb1: db 80h
120     304a 04          db  readf
121     304b 18          db  bdos1       ;sectors to read on track 1
122     304c 01          db  1           ;track 1
123     304d 01          db  1           ;sector 1
124     304e 800c       dw  cmpb+bdos0*128 ;base of second read
125     ;
126     3050          end

```

base	0078	33#	34	35	36	38	39	40
bdos	0806	18#						
bdoso	0019	28#	29	113	124			
bdos1	0018	29#	121					
bdose	1880	19#	25					
bdosl	1880	25#	27					
bdoss	0031	27#	29					
bias	0000	12#	15#	17	18	19	20	
boot	1600	20#	21	108				
bsw	00ff	41#	53					
coldstart	3007	52#	55					
cpmb	0000	17#	25	116	124			
dstat	0078	38#	70					
false	0000	7#	8	9				
ihigh	007a	40#	69					
ilow	0079	39#	67					
iopbo	3042	61	111#	117				
iopb1	3049	119#						
iopbl	0007	101	117#					
mon80	f800	31#						
ntrks	0002	26#	60					
rboot	1603	21#						
rbyte	007b	35#	50	86				
readf	0004	43#	112	120				
recal	0003	42#						
reset	007f	36#	57					
rmon80	ff0f	32#	80	89	94			
rstart	3000	46#	83	97				
rtype	0079	34#	49	75				
stack	0100	44#	47					
start	3015	63#	104					
testing	0000	9#	11	14	79	82	93	96
true	ffff	8#						
waito	301b	70#	72					



Appendix E: A Skeletal Cold Start Loader

```
; this is a sample cold start loader, which, when
; modified
; resides on track 00, sector 01 (the first sector on the
; diskette). we assume that the controller has loaded
; this sector into memory upon system start-up (this
; program can be keyed-in, or can exist in read/only
; memory
; beyond the address space of the cp/m version you are
; running). the cold start loader brings the cp/m system
; into memory at "loadp" (3400h + "bias"). in a 20k
; memory system, the value of "bias" is 0000h, with
; large
; values for increased memory sizes (see section 2).
; after
; loading the cp/m system, the cold start loader
; branches
; to the "boot" entry point of the bios, which begins at
; "bios" + "bias." the cold start loader is not used un-
; til the system is powered up again, as long as the bios
; is not overwritten. the origin is assumed at 0000h, an
; must be changed if the controller brings the cold start
; loader into another area, or if a read/only memory
; area
; is used.
```

```
0000                org 0                ; base of ram in
; cp/m

0014 =             msize equ 20         ; min mem size in
; kbytes
```

```

0000 =          bias      equ (msize-20)*1024    ; offset from 20k
                                           ; system
3400 =          ccp       equ 3400h+bias        ; base of the ccp
4a00 =          bios      equ ccp+1600h         ; base of the bios
0300 =          biosl     equ 0300h             ; length of the bios
4a00 =          boot      equ bios              ;
1900 =          size      equ bios+biosl-ccp     ; size of cp/m
                                           ; system
0032 =          sects     equ size/128          ; # of sectors to load

;          begin the load operation

cold:
0000 010200     lxi  b,2                        ; b=0, c=sector 2
0003 1632       mvi  d,sects                    ; d=# sectors to
                                           ; load
0005 210034     lxi  h,ccp                      ; base transfer
                                           ; address

lsect:          ; load the next sector

;          insert inline code at this point to
;          read one 128 byte sector from the
;          track given in register b, sector
;          given in register c,
;          into the address given by <hl>
; branch to location "cold" if a read error occurs
;
;
;          user supplied read operation goes
;          here...
;
;
0008 c36b00     jmp  past$patch                    ; remove this
                                           ; when patched
000b           ds    60h

past$patch:
; go to next sector if load is incomplete
006b 15        dcr  d                            ; sects=sects-1
006c ca004a    jz   boot                        ; head for the bios

;          more sectors to load
;
; we aren't using a stack, so use <sp> as scratch
; register
;          to hold the load address increment

006f 318000     lxi  sp,128                      ; 128 bytes per
                                           ; sector
0072 39        dad  sp                          ; <hl> = <hl> +
                                           128

```

```

0073 0c          inr  c          ; sector = sector + 1
0074 79          mov  a,c
0075 felb        cpi  27          ; last sector of
                                ; track?
0077 da0800      jc   lsect      ; no, go read
                                ; another

                                ; end of track, increment to next track

007a 0e01        mvi  c,l          ; sector = 1
007c 04          inr  b          ; track = track + 1
007d c30800      jmp  lsect        ; for another group
0080             end          ; of boot loader

```



Appendix F: CP/M Disk Definition Library

```
1: ;          CP/M 2.0 disk re-definition library
2: ;
3: ;          Copyright © 1979
4: ;          Digital Research
5: ;          Box 579
6: ;          Pacific Grove, CA
7: ;          93950
8: ;
9: ;          CP/M logical disk drives are defined using the
10: ;         macros given below, where the sequence of calls
11: ;         is:
12: ;
13: ;         disks  n
14: ;         diskdef parameter-list-0
15: ;         diskdef parameter-list-1
16: ;         ...
17: ;         diskdef parameter-list-n
18: ;         endef
19: ;
20: ;         where n is the number of logical disk drives attached
21: ;         to the CP/M system, and parameter-list-i defines the
22: ;         characteristics of the ith drive (i=0,1,...,n-1)
23: ;
24: ;         each parameter-list-i takes the form
25: ;             dn,fsc,lsc,[skf],bls,dks,dir,cks,ofs,[0]
26: ;         where
27: ;         dn      is the disk number 0,1,...,n-1
28: ;         fsc     is the first sector number (usually 0 or 1)
29: ;         lsc     is the last sector number on a track
30: ;         skf     is optional "skew factor" for sector translate
31: ;         bls     is the data block size (1024,2048,...,16384)
```

```

32: ;          dks      is the disk size in bls increments (word)
33: ;          dir      is the number of directory elements (word)
34: ;          cks      is the number of dir elements to checksum
35: ;          ofs      is the number of tracks to skip (word)
36: ;          [0]      is an optional 0 which forces 16K/directory end
37: ;
38: ;          for convenience, the form
39: ;              dn,dm
40: ;          defines disk dn as having the same characteristics as
41: ;          a previously defined disk dm.
42: ;
43: ;          a standard four drive CP/M system is defined by
44: ;              disks          4
45: ;              diskdef       0,1,26,6,1024,243,64,64,2
46: ;          dsk      set          0
47: ;              rept         3
48: ;          dsk      set          dsk+1
49: ;              diskdef      %dsk,0
50: ;              endm
51: ;              endef
52: ;
53: ;          the value of "begdat" at the end of assembly defines the
54: ;          beginning of the uninitialized ram area above the bios,
55: ;          while the value of "enddat" defines the next location
56: ;          following the end of the data area. the size of this
57: ;          area is given by the value of "datsiz" at the end of the
58: ;          assembly. note that the allocation vector will be quite
59: ;          large if a large disk size is defined with a small block
60: ;          size.
61: ;
62: dskhdr      macro    dn
63: ;;          define a single disk header list
64: dpe&dn:     dw        xlt&dn,0000h          ;translate table

```

```

65:          dw      0000h,0000h      ;scratch area
66:          dw      dirbuf,dpb&dn    ;dir buff,param block
67:          dw      csv&dn,alv&dn    ;check, alloc vectors
68:          endm
69: ;
70: disks          macro   nd
71: ;;            define nd disks
72: ndisks        set     nd          ;; for later reference
73: dpbase        equ     $           ;base of disk parameter blocks
74: ;;            generate the nd elements
75: dsknxt        set     0
76:              rept     nd
77:              dskhdr   %dsknxt
78: dsknxt        set     dsknxc+1
79:              endm
80:              endm
81: ;
82: dpbhdr          macro   dn
83: dpb&dn        equ     $           ;disk parm block
84:              endm
85: ;
86: ddb            macro   data,comment
87: ;;            define a db statement
88:              db      data          comment
89:              endm
90: ;
91: ddw            macro   data,comment
92: ;;            define a dw statement
93:              dw      data          comment
94:              endm
95: ;
96: gcd            macro   m,n
97: ;;            greatest common divisor of m,n
98: ;;            produces value gcdn as result

```

```

99: ;;                                (used in sector translate table generation)
100: gcdm                             set      m                ;;variable for m
101: gcdn                             set      n                ;;variable for n
102: gcdr                             set      0                ;;variable for r
103:                                 rept      65535
104: gcdx                             set      gcdm/gcdn
105: gcdr                             set      gcdm-gcdx*gcdn
106:                                 if        gcdr = 0
107:                                 exitm
108:                                 endif
109: gcdm                             set      gcdn
110: gcdn                             set      gcdr
111:                                 endm
112:                                 endm
113: ;
114: diskdef                           macro   dn,fsc,lsc,skf,bls,dks,dir,cks,ofs,k16
115: ;;                                generate the set statements for later tables
116:                                 if      nul lsc
117: ;;                                current disk dn                same as previous fsc
118: dpb&dn                             equ     dpb&fsc           ;;equivalent parameters
119: als&dn                             equ     als&fsc           ;;same allocation vector size
120: css&dn                             equ     css&fsc           ;;same checksum vector size
121: xlt&dn                             equ     xlt&fsc           ;;same translate table
122:                                 else
123: secmax                             set     lsc-(fsc)         ;;sectors 0...secmax
124: sectors                             set     secmax+1         ;;number of sectors
125: als&dn                             set     (dks)/8          ;;size of allocation vector
126:                                 if      ((dks) mod 8) ne 0
127: als&dn                             set     als&dn+1
128:                                 endif
129: css&dn                             set     (cks)/4          ;;number of checksum elements
130: ;;                                generate the block shift value
131: blkval                             set     bls/128          ;;number of sectors/ block

```

```

132: blkshf          set      0          ;;counts right 0's in blkval
133: blkmsk         set      0          ;;fills with l's from right
134:                rept     16         ;;once for each bit position
135:                if      blkval=1
136:                exitm
137:                endif
138: ;;              otherwise, high order 1 not found yet
139: blkshf          set      blkshf+1
140: blkmsk         set      (blkmsk shl l) or l
141: blkval          set      blkval/2
142:                endm
143: ;;              generate the extent mask byte
144: blkval          set      bls/1024        ;;number of kilobytes/ block
145: extmsk         set      0          ;;fill from right with l's
146:                rept     16
147:                if      blkval=1
148:                exitm
149:                endif
150: ;;              otherwise more to shift
151: extmsk         set      (extmsk shl l) or l
152: blkval          set      blkval/2
153:                endm
154: ;;              may be double byte allocation
155:                if      (dks) > 256
156: extmsk         set      (extmsk shr l)
157:                endif
158: ;;              may be optional [0] in last position
159:                if      not nul k16
160: extmsk         set      k16
161:                endif
162: ;;              now generate directory reservation bit vector
163: dirrem         set      dir          ;;# remaining to process
164: dirbks         set      bls/32        ;;number of entries per block
165: dirblk         set      0          ;;fill with l's on each loop
166:                rept     16
167:                if      dirrem=0
168:                exitm
169:                endif

```

```

170: ;;          not complete, iterate once again
171: ;;          shift right and add 1 high order bit
172: dirblk      set          (dirblk shr l) or 8000h
173:             if          dirrem > dirbks
174: dirrem      set          dirrem-dirbks
175:             else
176: direem      set          0
177:             endif
178:             endm
179:             dpbhdr   dn          ;;generate equ $
180:             ddw      %sectors,<;sec per track>
181:             ddb      %blkshf,<;block shift>
182:             ddb      %blkmsk,<;block mask>
183:             ddb      %extmsk,<;extnt mask>
184:             ddw      %(dks)-1,<;disk size-1>
185:             ddw      %(dir)-1,<;directory max>
186:             ddb      %dirblk shr 8,<;alloc0>
187:             ddb      %dirblk and 0ffh,<;alloc1>
188:             ddw      %(cks)/4,<;check size>
189:             ddw      %ofs,<;offset>
190: ;;          generate the translate table, if requested
191:             if          nul skf
192: xlt&dn      equ          0          ;no xlate table
193:             else
194:             if          skf = 0
195: xlt&dn      equ          0          ;no xlate table
196:             else
197: ;;          generate the translate table
198: nxtsec      set          0          ;;next sector to fill
199: nxtbas      set          0          ;;moves by one on overflow
200:             gcd      %sectors,skf
201: ;;          gcdn = gcd(sectors,skew)
202: neltst      set          sectors/gcdn
203: ;;          neltst is number of elements to generate

```

```

204: ;; before we overlap previous elements
205: nelts set nelts ;;counter
206: xlt&dn equ $ ;translate table
207: rept sectors ;;once for each sector
208: if sectors < 256
209: ddb %nxtsec+(fsc)
210: else
211: ddw %nxtsec+(fsc)
212: endif
213: nxtsec set nxtsec+(skf)
214: if nxtsec >= sectors
215: set nxtsec-sectors
216: endif
217: nelts set nelts-1
218: if nelts = 0
219: nxtbas set nxtbas+1
220: nxtsec set nxtbas
221: nelts set nelts
222: endif
223: endm
224: endif ;;end of nul fac test
225: endif ;;end of nul bls test
226: endm
227: ;
228: defds macro lab,space
229: lab: ds space
230: endm
231: ;
232: lds macro lb,dn,val
233: defds lb&dn,%val&dn
234: endm
235: ;
236: endef macro
237: ;; generate the necessary ram data areas

```

```
238: begdat      equ      $
239: dirbuf:     ds       128          ;directory access buffer
240: dsknxt      set      0
241:             rept     ndisks      ;;once for each disk
242:             lds      alv,%dsknxt,als
243:             lds      csv,%dsknxt,ccs
244: dsknxt      set      dsknxt+1
245:             endm
246: enddat      equ      $
247: datsiz      equ      $-begdat
248: ;;          db 0 at this point forces hex record
249:             endm
```

Appendix G: Blocking and Deblocking Algorithms

```

1      ;
2      ;
3      ;           sector deblocking algorithms for cp/m 2.0
4      ;
5      ;
6      ;
7      ;           utility macro to compute sector mask
8      smask macro   hblk
9      ;;         compute log2(hblk), return @x as result
10     ;;         (2 ** @x = hblk on return)
11     @y set      hblk
12     @x set      0
13     ;;         count right shifts of @y until = 1
14     rept      8
15     if        @y = 1
16     exitm
17     endif
18     ;;         @y is not 1, shift right one position
19     @y set     @y shr 1
20     @x set     @x + 1
21     endm
22     endm
23     ;
24     ;
25     ;
26     ;           cp/m to host disk constants
27     ;
28     ;
29     0800 = blksiz equ 2048 ;cp/m allocation size
30     0200 = hstsiz equ 512 ;host disk sector size
31     0014 = hstspt equ 20 ;host disk sectors/trk
32     0004 = hstblk equ hstsiz/128 ;cp/m sects/host buff
33     0050 = cpmspt equ hstblk * hstspt ;cp/m sectors/track

```

```

34      0003 =   secmsk   equ      hstblk-1           ;sector mask
35                                     smask   hstblk           ;compute sector mask
36      0002 =   secshf   equ      @x             ;log2(hstblk)
37      ;
38      ;
39      ;
40      ;           bdos constants on entry to write
41      ;
42      ;
43      0000 =   wrall    equ      0              ;write to allocated
44      0001 =   wrdir    equ      1              ;write to directory
45      0002 =   wrual    equ      2              ;write to unallocated
46      ;
47      ;
48      ;
49      ;           the bdos entry points given below show the
50      ;           code which is relevant to deblocking only.
51      ;
52      ;
53      ;
54      ;           diskdef macro, or hand coded tables go here
55      0000 =   dpbase   equ      $              ;disk param block base
56      ;
57      boot:
58      wboot:
59      ;enter here on system boot to initialize
60      0000 af         xra      a                ;0 to accumulator
61      0001 326a01     sta      hstact           ;host buffer inactive
62      0004 326c01     sta      unacnt          ;clear unalloc count
63      0007 c9         ret
64      ;
65      home:
66      ;home the selected disk
67      home:
68      0008 3a6b01     lda      hstwrnt         ;check for pending write
69      000b b7         ora      a
70      000c c21200     jnz      homed
71      000f 326a01     sta      hstact           ;clear host active flag
72      homed:
73      0012 c9         ret
74      ;
75      seldsk:
76      ;select disk
77      0013 79         mov      a,c                ;selected disk number
78      0014 326101     sta      sekdisk           ;seek disk number
79      0017 6f         mov      l,a                ;disk number to hl
80      0018 2600     mvi      h,0
81      rept      4           ;multiply by 16
82      dad      h
83      endm
84      001a+29         dad      h
85      001b+29         dad      h
86      001c+29         dad      h
87      001d+29         dad      h
88      001e 110000     lxi      d,dpbase          ;base of parm block

```

```

89      0021 19      dad      d          ;hl=.dpb(curdsk)
90      0022 c9      ret
91      ;
92      ;settrk:
93      ;set track given by registers bc
94      0023 60      mov      h,b
95      0024 69      mov      l,c
96      0025 226201  shld    sektrk      ;track to seek
97      0028 c9      ret
98      ;
99      ;setsec:
100     ;set sector given by register c
101     0029 79      mov      a,c
102     002a 326401  sta      seksec      ;sector to seek
103     002d c9      ret
104     ;
105     ;setdma:
106     ;set dma address given by bc
107     002e 60      mov      h,b
108     002f 69      mov      l,c
109     0030 227501  shld    dmaadr
110     0033 c9      ret
111     ;
112     ;sectran:
113     ;translate sector number bc
114     0034 60      mov      h,b
115     0035 69      mov      l,c
116     0036 c9      ret
117     ;
118     ;
119     ;
120     ;          the read entry point takes the place of
121     ;          the previous bios definition for read.
122     ;
123     ;
124     ;          read:
125     ;read the selected cp/m sector
126     0037 af      xra      a
127     0038 326c01  sta      unacct
128     003b 3e01    mvi      a,1
129     003d 327301  sta      readop      ;read operation
130     0040 327201  sta      rsflag      ;must read data
131     0043 3e02    mvi      a,wruall
132     0045 327401  sta      wrtype      ;treat as unalloc
133     0048 c3b600  jmp      rwoper      ;to perform the read
134     ;
135     ;
136     ;
137     ;          the write entry point takes the place of
138     ;          the previous bios definition for write.
139     ;
140     ;
141     ;          write:
142     ;write the selected cp/m sector
143     004b af      xra      a          ;0 to accumulator

```

```

144      004c 327301      sta      readop      ;not a read operation
145      004f 79         mov      a,c         ;write type in c
146      0050 327401      sta      wrtype
147      0053 fe02       cpi      wrual      ;write unallocated?
148      0055 c26f00     jnz      chkuna     ;check for unalloc
149      ;
150      ;               write to unallocated, set parameters
151      0058 3e10       mvi      a,blksiz/128 ;next unalloc recs
152      005a 326c01     sta      unacnt
153      005d 3a6101     lda      sekdisk    ;disk to seek
154      0060 326d01     sta      unadsk     ;unadsk = sekdisk
155      0063 2a6201     lhld    sektrk
156      0066 226e01     shld    unatrkl    ;unatrkl = sektrk
157      0069 3a6401     lda      seksec
158      006c 327001     sta      unasec    ;unasec = seksec
159      ;
160      ;               chkuna:
161      ;               ;check for write to unallocated sector
162      006f 3a6c01     lda      unacnt    ;any unalloc remain?
163      0072 b7         ora      a
164      0073 caae00     jz      alloc      ;skip if not
165      ;
166      ;               more unallocated records remain
167      0076 3d         dcr      a         ;unacnt = unacnt-1
168      0077 326c01     sta      unacnt
169      007a 3a6101     lda      sekdisk    ;same disk?
170      007d 216d01     lxi      h,unadsk
171      0080 be         cmp      m         ;sekdisk = unadsk?
172      0081 c2ae00     jnz      alloc      ;skip if not
173      ;
174      ;               disks are the same
175      0084 216e01     lxi      h,unatrkl
176      0087 cd5301     call    sektrkcmp  ;sektrkl = unatrkl?
177      008a c2ae00     jnz      alloc      ;skip if not
178      ;
179      ;               tracks are the same
180      008d 3a6401     lda      seksec    ;same sector?
181      0090 217001     lxi      h,unasec
182      0093 be         cmp      m         ;seksec = unasec?
183      0094 c2ae00     jnz      alloc      ;skip if not
184      ;
185      ;               match, move to next sector for future ref
186      0097 34         inr      m         ;unasec = unasec+1
187      0098 7e         mov      a,m       ;end of track?
188      0099 fe50       cpi      cpmspt    ;count cp/m sectors
189      009b daa700     jc      noovf     ;skip if no overflow
190      ;
191      ;               overflow to next track
192      009e 3600       mvi      m,o       ;unasec = 0
193      00a0 2a6e01     lhld    unatrkl
194      00a3 23         inx      h
195      00a4 226e01     shld    unatrkl    ;unatrkl = unatrkl+1
196      ;
197      ;               noovf:
198      ;               ;match found, mark as unnecessary read
199      00a7 af         xra      a         ;0 to accumulator

```

```

200      00ab 327201      sta      rsflag          ;rsflag = 0
201      00ab c3b600      jmp      rwoper          ;to perform the write
202      ;
203      alloc:
204      ;not an unallocated record, requires pre-read
205      00ae af          xra      a              ;0 to accum
206      00af 326c01      sta      unacct         ;unacct = 0
207      00b2 3c          inr      a              ;1 to accum
208      00b3 327201      sta      rsflag = 1     ;rsflag = 1
209      ;
210      ;
211      ;
212      ;          common code for read and write follows
213      ;
214      ;
215      rwoper:
216      ;enter here to perform the read/write
217      00b6 af          xra      a              ;zero to accum
218      00b7 327101      sta      erflag         ;no errors (yet)
219      00ba 3a6401      lda      seksec         ;compute host sector
220      rept      secshf
221      ora      a              ;carry = 0
222      rar              ;shift right
223      endm
224      00bd+b7          ora      a              ;carry = 0
225      00be+1f          rar              ;shift right
226      00bf+b7          ora      a              ;carry = 0
227      00c0+1f          rar              ;shift right
228      00c1 326901      sta      sekhst         ;host sector to seek
229      ;
230      ;          active host sector?
231      00c4 216a01      lxi      h,hstact       ;host active flag
232      00c7 7e          mov      a,m
233      00c8 3601        mvi      m,1           ;always becomes 1
234      00ca b7          ora      a              ;was it already?
235      00cb caf200      jz      filhst         ;fill host if not
236      ;
237      ;          host buffer active, same as seek buffer?
238      00ce 3a6101      lda      sekdisk
239      00d1 216501      lxi      h,hstdsk      ;same disk?
240      00d4 be          cmp      m              ;sekdisk = hstdsk?
241      00d5 c2eb00      jnz      nomatch
242      ;
243      ;          same disk, same track?
244      00d8 216601      lxi      h,hsttrk
245      00db cd5301      call    sektrkcmp      ;sektrk = hsttrk?
246      00de c2eb00      jnz      nomatch
247      ;
248      ;          same disk, same track, same buffer?
249      00e1 3a6901      lda      sekhst
250      00e4 216801      lxi      h,hstsec      ;sekhst = hstsec?
251      00e7 be          cmp      m
252      00e8 ca0f01      jz      match          ;skip if match
253      ;
254      nomatch:

```

```

255                                     ;proper disk, but not correct sector
256 00eb 3a6b01 lda hstwrtn ;host written?
257 00ee b7 ora a
258 00ef c45f01 cnz writehst ;clear host buff
259 ;
260 filhst:
261                                     ;may have to fill the host buffer
262 00f2 3a6101 lda sekdisk
263 00f5 326501 sta hstdsk
264 00f8 2a6201 lhld sektrk
265 00fb 226601 shld hsttrk
266 00fe 3a6901 lda sekhst
267 0101 326801 sta hstsec
268 0104 3a7201 lda rsflag ;need to read?
269 0107 b7 ora a
270 0108 c46001 cnz readhst ;yes, if 1
271 010b af xra a ;0 to accum
272 010c 326b01 sta hstwrtn ;no pending write
273 ;
274 match:
275                                     ;copy data to or from buffer
276 010f 3a6401 lda seksec ;mask buffer number
277 0112 e603 ani secmsk ;least signif bits
278 0114 6f mov l,a ;ready to shift
279 0115 2600 mvi h,0 ;double count
280 rept 7 ;shift left 7
281 dad h
282 endm
283 0117+29 dad h
284 0118+29 dad h
285 0119+29 dad h
286 011a+29 dad h
287 011b+29 dad h
288 011c+29 dad h
289 011d+29 dad h
290 ; hl has relative host buffer address
291 011e 117701 lxi d,hstbuf
292 0121 19 dad d ;hl = host address
293 0122 eb xchg ;now in de
294 0123 2a7501 lhld dmaadr ;get/put cp/m data
295 0126 0e80 mvi c,128 ;length of move
296 0128 3a7301 lda readop ;which way?
297 012b b7 ora a
298 012c c23501 jnz rwmov ;skip if read
299 ;
300 ; write operation, mark and switch direction
301 012f 3e01 mvi a,1
302 0131 326b01 sta hstwrtn ;hstwrtn = 1
303 0134 eb xchg ;source/dest swap
304 ;
305 rwmov:
306                                     ;c initially 128, de is source, hl is dest
307 0135 1a ldax d ;source character
308 0136 13 inx d
309 0137 77 mov m,a ;to dest

```

```

310      0138 23          inx      h
311      0139 od          dcr      c                ;loop 128 times
312      013a c23501     jnz      rwmove
313      ;
314      ;
315      013d 3a7401     lda      wrtype          ;write type
316      0140 fe01       cpi      wrdir           ;to directory?
317      0142 3a7101     lda      erflag          ;in case of errors
318      0145 c0         rnz      ;no further processing
319      ;
320      ;
321      0146 b7         ora      a                ;errors?
322      0147 c0         rnz      ;skip if so
323      0148 af         xra      a                ;0 to accum
324      0149 326b01     sta      hstwr           ;buffer written
325      014c cd5f01     call    writehst
326      014f 3a7101     lda      erflag
327      0152 c9         ret
328      ;
329      ;
330      ;
331      ;
332      ;
333      ;
334      ;
335      ;
336      ;
337      ;
338      ;
339      ;
340      ;
341      ;
342      ;
343      ;
344      ;
345      ;
346      ;
347      ;
348      ;
349      ;
350      ;
351      ;
352      ;
353      ;
354      ;
355      ;
356      ;
357      ;
358      ;
359      ;
360      ;
361      ;
362      ;
363      ;
364      ;

```

```

365                                     ;into hstbuf and return error flag in erflag.
366 0160 c9                             ret
367                                     ;
368                                     ;
369                                     ;
370                                     ; uninitialized ram data areas
371                                     ;
372                                     ;
373                                     ;
374 0161 sekdisk: ds 1                   ;seek disk number
375 0162 sektrk: ds 2                   ;seek track number
376 0164 seksec: ds 1                   ;seek sector number
377                                     ;
378 0165 hstdsk: ds 1                   ;host disk number
379 0166 hsttrk: ds 2                   ;host track number
380 0168 hstsec: ds 1                   ;host sector number
381                                     ;
382 0169 sekhst: ds 1                   ;seek shr secshf
383 016a hstact: ds 1                   ;host active flag
384 016b hstwrtr: ds 1                  ;host written flag
385                                     ;
386 016c unacct: ds 1                   ;unalloc rec cnt
387 016d unadsk: ds 1                   ;last unalloc disk
388 016e unatrk: ds 2                   ;last unalloc track
389 0170 unasec: ds 1                   ;last unalloc sector
390                                     ;
391 0171 erflag: ds 1                    ;error reporting
392 0172 rsflag: ds 1                    ;read sector flag
393 0173 readop: ds 1                    ;1 if read operation
394 0174 wrtype: ds 1                    ;write operation type
395 0175 dmaadr: ds 2                    ;last dma address
396 0177 hstbuf: ds hstsiz               ;host buffer
397                                     ;
398                                     ;
399                                     ;
400                                     ; the endef macro invocation goes here
401                                     ;
402                                     ;
403 0377 end

```

alloc	00ae	164	172	177	183	203#		
blksiz	0800	29#	151					
boot	0000	57#						
chkuna	006f	148	160#					
cpmspt	0050	33#	188					
dmaadr	0175	109	294	395#				
dpbase	0000	55#	88					
erflag	0171	218	317	326	391#			
filhst	00f2	235	260#					
home	0008	65#	67#					
homed	0012	70	72#					
hstact	016a	61	71	231	383#			
hstblk	0004	32#	33	34	35			
hstbuf	0177	291	396#					
hstdsk	0165	239	263	378#				
hstsec	0168	250	267	380#				
hstsiz	0200	30#	32	396				
hstspt	0014	31#	33					
hstrk	0166	244	265	379#				
hstwrt	016b	68	256	272	302	324	384#	
match	010f	252	274#					
nomatch	00eb	241	246	254#				
noovf	00a7	189	197#					
read	0037	124#						
readhst	0160	270	362#					
readop	0173	129	144	296	393#			
rsflag	0172	130	200	208	268	392#		
rwmove	0135	298	305#	312				
rwoper	00b6	133	201	215#				
secmsk	0003	34#	277					
secshf	0002	36#	220					
sectran	0034	112#						
sekdsd	0161	78	153	169	238	262	374#	
sekhst	0169	228	249	266	382#			
seksec	0164	102	157	180	219	276	376#	
sektrk	0162	96	155	264	337	375#		
sektrkcmp	0153	176	245	334#				
seldsk	0013	75#						
setdma	002e	105#						
setsec	0029	99#						
settrk	0023	92#						
unacnt	016c	62	127	152	162	168	206	386#
unadsk	016d	154	170	387#				
unasec	0170	158	181	389#				
unatrk	016e	156	175	193	195	388#		
wboot	0000	58#						
wrall	0000	43#						
wrdir	0001	44#	316					
write	004b	141#						
writehst	015f	258	325	355#				
wrtype	0174	132	146	315	394#			
wruval	0002	45#	131	147				

17	11	Search For First	DE = FCB Address	A = Directory Code
18	12	Search For Next	none	A = Directory Code
19	13	Delete File	DE = FCB Address	A = none
20	14	Read Sequential	DE = FCB Address	A = Error Code
21	15	Write Sequential	DE = FCB Address	A = Error Code
22	16	Make File	DE = FCB Address	A = FF if no DIR Space
23	17	Rename File	DE = FCB Address	A = FF if not found
24	18	Return Login Vector	none	HL = Login Vector*
25	19	Return Current Disk	none	A = Current Disk Number
26	1A	Set DMA Address	DE = DMA Address	none
27	1B	Get ADDR (ALLOC)	none	HL = ALLOC Address*
28	1C	Write Protect Disk	none	none
29	1D	Get Read/only Vector	none	HL = R/O Vector Value*
30	1E	Set File Attributes	DE = FCB Address	A = none
31	1F	Get ADDR (Disk Parms)	none	HL = DPB Address
32	20	Set/Get User Code	E = 0FFH for Get E = 00 to 0FH for Set	User Number
33	21	Read Random	DE = FCB Address	A = Error Code
34	22	Write Random	DE = FCB Address	A = Error Code
35	23	Compute File Size	DE = FCB Address	r0, r1, r2
36	24	Set Random Record	DE = FCB Address	r0, r1, r2
37	25	Reset Drive	DE = Drive Vector	A = 0
38	26	Access Drive	not supported	
39	27	Free Drive	not supported	
40	28	Write Random with Fill	DE = FCB	A = Error Code

*Note that A = L, and B = H upon return.